

LaTeX2 **Functional** Interfaces for LaTeX3 Programming Layer

Jianrui Lyu (tolvjr@163.com)
<https://github.com/lvjrr/functional>

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LaTeX3 programming layer (**expl3**) is very powerful for advanced users, but it is a little complicated for normal users. This **functional** package aims to provide intuitive LaTeX2 functional interfaces for it.

Although there are functions in LaTeX3, the evaluation of them is from outside to inside. With this package, the evaluation of functions is from inside to outside, which is the same as other programming languages such as **JavaScript** or **Lua**. In this way, it is rather easy to debug code too.

Note that many paragraphs in this manual are copied from the documentation of **expl3**.

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Chapter 1

Overview of Features

1.1 Evaluation from Inside to Outside

We will compare our first example with a similar Lua example:

<pre>\IgnoreSpacesOn \PrgNewFunction \MathSquare { m } { \IntSet \lTmpaInt { \IntEval {#1 * #1} } \Result { \Value \lTmpaInt } } \IgnoreSpacesOff \MathSquare{5} \MathSquare{\MathSquare{5}}</pre>	<pre>-- define a function -- function MathSquare (arg) local lTmpaInt = arg * arg return lTmpaInt end -- use the function -- print(MathSquare(5)) print(MathSquare(MathSquare(5)))</pre>
--	--

Both examples calculate first the square of 5 and produce 25, then calculate the square of 25 and produce 625. In contrast to `expl3`, this `functional` package does evaluation of functions from inside to outside, which means composition of functions works like other programming languages such as Lua or JavaScript.

You can define new functions with `\PrgNewFunction` command. To make composition of functions work as expected, every function *must not* insert directly any token to the input stream. Instead, a function *must* pass the result (if any) to `functional` package with `\Result` command. And `functional` package is responsible for inserting result tokens to the input stream at the appropriate time.

To remove space tokens inside function code in defining functions, you'd better put function definitions inside `\IgnoreSpacesOn` and `\IgnoreSpacesOff` block. Within this block, `~` is used to input a space.

At the end of this section, we will compare our factorial example with a similar Lua example:

<pre>\IgnoreSpacesOn \PrgNewFunction \Fact { m } { \IntCompareTF {#1} = {0} { \Result {1} }{ \Result {\IntMathMult{#1}{\Fact{\IntMathSub{#1}{1}}}} } } \IgnoreSpacesOff \Fact{4}</pre>	<pre>-- define a function -- function Fact (n) if n == 0 then return 1 else return n * Fact(n-1) end end -- use the function -- print(Fact(4))</pre>
--	--

1.2 Group Scoping of Functions

In Lua language, a function or a condition expression makes a block, and the values of local variables will be reset after a block. In `functional` package, a condition expression is in fact a function, and you can make every function become a group by setting `\Functional{scoping=true}`. For example

<code>\Functional{scoping=true}</code>		<code>-- lua code --</code>
<code>\IgnoreSpacesOn</code>		<code>-- begin example --</code>
<code>\IntSet \lTmpaInt {1}</code>		<code>local a = 1</code>
<code>\IntLogVar \lTmpaInt</code>	<code>% ---- 1</code>	<code>print(a) ---- 1</code>
<code>\PrgNewFunction \SomeFun { } {</code>		<code>function SomeFun()</code>
<code>\IntSet \lTmpaInt {2}</code>		<code>local a = 2</code>
<code>\IntLogVar \lTmpaInt</code>	<code>% ---- 2</code>	<code>print(a) ---- 2</code>
<code>\IntCompareTF {1} > {0} {</code>		<code>if 1 > 0 then</code>
<code>\IntSet \lTmpaInt {3}</code>		<code>local a = 3</code>
<code>\IntLogVar \lTmpaInt</code>	<code>% ---- 3</code>	<code>print(a) ---- 3</code>
<code>}{ }</code>		<code>end</code>
<code>\IntLogVar \lTmpaInt</code>	<code>% ---- 2</code>	<code>print(a) ---- 2</code>
<code>}</code>		<code>end</code>
<code>\SomeFun</code>		<code>SomeFun()</code>
<code>\IntLogVar \lTmpaInt</code>	<code>% ---- 1</code>	<code>print(a) ---- 1</code>
<code>\IgnoreSpacesOff</code>		<code>-- end example --</code>

Same as `expl3`, the names of local variables *must* start with `l`, while names of global variables *must* start with `g`. The difference is that `functional` package provides only one function for setting both local and global variables of the same type, by checking leading letters of their names. So for integer variables, you can write `\IntSet\lTmpaInt{1}` and `\IntSet\gTmbpInt{2}`.

The previous example will produce different result if we change variable from `\lTmpaInt` to `\gTmpaInt`.

<code>\Functional{scoping=true}</code>		<code>-- lua code --</code>
<code>\IgnoreSpacesOn</code>		<code>-- begin example --</code>
<code>\IntSet \gTmpaInt {1}</code>		<code>a = 1</code>
<code>\IntLogVar \gTmpaInt</code>	<code>% ---- 1</code>	<code>print(a) ---- 1</code>
<code>\PrgNewFunction \SomeFun { } {</code>		<code>function SomeFun()</code>
<code>\IntSet \gTmpaInt {2}</code>		<code>a = 2</code>
<code>\IntLogVar \gTmpaInt</code>	<code>% ---- 2</code>	<code>print(a) ---- 2</code>
<code>\IntCompareTF {1} > {0} {</code>		<code>if 1 > 0 then</code>
<code>\IntSet \gTmpaInt {3}</code>		<code>a = 3</code>
<code>\IntLogVar \gTmpaInt</code>	<code>% ---- 3</code>	<code>print(a) ---- 3</code>
<code>}{ }</code>		<code>end</code>
<code>\IntLogVar \gTmpaInt</code>	<code>% ---- 3</code>	<code>print(a) ---- 3</code>
<code>}</code>		<code>end</code>
<code>\SomeFun</code>		<code>SomeFun()</code>
<code>\IntLogVar \gTmpaInt</code>	<code>% ---- 3</code>	<code>print(a) ---- 3</code>
<code>\IgnoreSpacesOff</code>		<code>-- end example --</code>

As you can see, the values of global variables will never be reset after a group.

1.3 Tracing Evaluation of Functions

Since every function in `functional` package will pass its return value to the package, it is quite easy to debug your code. You can turn on the tracing by setting `\Functional{tracing=true}`. For example, the tracing log of the first example in this chapter will be the following:

```

[I] \MathSquare{5}
    [I] \IntEval{5*5}
        [I] \Expand{\int_eval:n {5*5}}
        [0] 25
    [I] \Result{25}
    [0] 25
[0] 25
[I] \IntSet{\lTmptInt }{25}
[0]
    [I] \Value{\lTmptInt }
    [0] 25
[I] \Result{25}
[0] 25
[0] 25
[I] \MathSquare{25}
    [I] \IntEval{25*25}
        [I] \Expand{\int_eval:n {25*25}}
        [0] 625
    [I] \Result{625}
    [0] 625
[0] 625
[I] \IntSet{\lTmptInt }{625}
[0]
    [I] \Value{\lTmptInt }
    [0] 625
[I] \Result{625}
[0] 625
[0] 625

```

1.4 Definitions of Functions

Within `expl3`, there are eight commands for defining new functions, which is good for power users.

<code>\cs_new:Npn</code>	<code>\cs_new:Nn</code>
<code>\cs_new_nopar:Npn</code>	<code>\cs_new_nopar:Nn</code>
<code>\cs_new_protected:Npn</code>	<code>\cs_new_protected:Nn</code>
<code>\cs_new_protected_nopar:Npn</code>	<code>\cs_new_protected_nopar:Nn</code>

Within `functional` package, there is only one command (`\PrgNewFunction`) for defining new functions, which is good for normal users. The created functions are always protected and accept `\par` in their arguments.

Since `functional` package gets the results of functions by evaluation (including expansion and execution by \TeX), it is natural to protect all functions.

1.5 Variants of Arguments

Within `expl3`, there are several expansion variants for arguments, and many expansion functions for expanding them, which are necessary for power users.

<code>\module_foo:c</code>	<code>\exp_args:Nc</code>
<code>\module_bar:e</code>	<code>\exp_args:Ne</code>
<code>\module_bar:x</code>	<code>\exp_args:Nx</code>
<code>\module_bar:f</code>	<code>\exp_args:Nf</code>
<code>\module_bar:o</code>	<code>\exp_args:No</code>
<code>\module_bar:V</code>	<code>\exp_args:NV</code>
<code>\module_bar:v</code>	<code>\exp_args:Nv</code>

Within `functional` package, there are only three variants (`c`, `e`, `V`) are provided, and these variants are defined as functions (`\Name`, `\Expand`, `\Value`, respectively), which are easier to use for normal users.

```
\newcommand\test{uvw}  
\Name{test}
```

uvw

```
\newcommand\test{uvw}  
\Expand{111\test222}
```

111uvw222

```
\IntSet\lTmptInt{123}  
\Value\lTmptInt
```

123

The most interesting feature is that you can compose these functions. For example, you can easily get the `v` variant of `expl3` by simply composing `\Name` and `\Value` functions:

```
\IntSet\lTmptInt{123}  
\Value{\Name{\lTmptInt}}
```

123

Chapter 2

Functional Programming (Prg)

2.1 Defining Functions and Conditionals

\PrgNewFunction $\langle function \rangle$ $\{\langle argument specification \rangle\}$ $\{\langle code \rangle\}$

Creates protected $\langle function \rangle$ for evaluating the $\langle code \rangle$. Within the $\langle code \rangle$, the parameters ($\#1$, $\#2$, *etc.*) will be replaced by those absorbed by the function. The returned value must be passed with **\Result** function. The definition is global and an error results if the $\langle function \rangle$ is already defined.

The $\{\langle argument specification \rangle\}$ in a list of letters, where each letter is one of the following argument specifiers (nearly all of them are M or m for functions provided by this package):

- M single-token argument, which will be manipulated first
- m multi-token argument, which will be manipulated first
- N single-token argument, which will not be manipulated first
- n multi-token argument, which will not be manipulated first

The argument manipulation for argument type M or m is: if the argument starts with a function defined with **\PrgNewFunction**, the argument will be evaluated and replaced with the returned value.

\PrgNewConditional $\langle function \rangle$ $\{\langle argument specification \rangle\}$ $\{\langle code \rangle\}$

Creates protected conditional $\langle function \rangle$ for evaluating the $\langle code \rangle$. The returned value of the $\langle function \rangle$ must be either **\cTrueBool** or **\cFalseBool** and be passed with **\Result** function.. The definition is global and an error results if the $\langle function \rangle$ is already defined.

Assume the $\langle function \rangle$ is **\FooIfBar**, then another function **\FooIfBarTF** will be created at the same time. **\FooIfBarTF** function has two extra arguments which are $\{\langle true code \rangle\}$ and $\{\langle false code \rangle\}$.

2.2 Collecting Returned Values

\Result $\{\langle tokens \rangle\}$

Appends $\langle tokens \rangle$ to **\gResultTl**, which holds the returned value of current function. This function is normally used in the $\langle code \rangle$ of **\PrgNewFunction** and **\PrgNewConditional**.

Chapter 3

Argument Using (Use)

3.1 Expanding Tokens

\Name {*<control sequence name>*}

Expands the *<control sequence name>* until only characters remain, then converts this into a control sequence and returns it. The *<control sequence name>* must consist of character tokens when exhaustively expanded.

\Value *<variable>*

Recovers the content of a *<variable>* and returns the value. An error is raised if the variable does not exist or if it is invalid. Note that it is the same as **\TlUse** for *<tl var>*, or **\IntUse** for *<int var>*.

\Expand {*<tokens>*}

Expands the *<tokens>* exhaustively and returns the result.

\ExpNot {*<tokens>*}

Prevents expansion of the *<tokens>* inside the argument of **\Expand** function. The argument of **\ExpNot** must be surrounded by braces.

\ExpValue *<variable>*

Recovers the content of the *<variable>*, then prevents expansion of this material inside the argument of **\Expand** function.

3.2 Using Tokens

\UseOne {*<argument>*}

\GobbleOne {*<argument>*}

The function **\UseOne** absorbs one argument and returns it. **\GobbleOne** absorbs one argument and returns nothing. For example

```
\UseOne{abc}\GobbleOne{ijk}\UseOne{xyz}
```

```
abcxyz
```

```
\UseGobble {⟨arg1⟩} {⟨arg2⟩}  
\GobbleUse {⟨arg1⟩} {⟨arg2⟩}
```

These functions absorb two arguments. The function `\UseGobble` discards the second argument, and returns the content of the first argument. `\GobbleUse` discards the first argument, and returns the content of the second argument. For example

```
\UseGobble{abc}{uvw}\GobbleUse{abc}{uvw}
```

```
abcuvw
```

Chapter 4

Control Structures (Bool)

4.1 Constant and Scratch Booleans

`\cTrueBool \cFalseBool`

Constants that represent `true` and `false`, respectively. Used to implement predicates. For example

```
\BoolVarIfTF \cTrueBool {\Result{True!}} {\Result{False!}}
\BoolVarIfTF \cFalseBool {\Result{True!}} {\Result{False!}}
```

True! False!

`\lTmptaBool \lTmptbBool \lTmptcBool \lTmptiBool \lTmptjBool \lTmptkBool`

Scratch booleans for local assignment. These are never used by the `functional` package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

`\gTmptaBool \gTmptbBool \gTmptcBool \gTmptiBool \gTmptjBool \gTmptkBool`

Scratch booleans for global assignment. These are never used by the `functional` package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

4.2 Creating and Setting Booleans

`\BoolNew` $\langle boolean \rangle$

Creates a new $\langle boolean \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle boolean \rangle$ is initially `false`.

`\BoolConst` $\langle boolean \rangle$ $\{\langle boolexpr \rangle\}$

Creates a new constant $\langle boolean \rangle$ or raises an error if the name is already taken. The value of the $\langle boolean \rangle$ is set globally to the result of evaluating the $\langle boolexpr \rangle$. For example

```
\BoolConst \cFooSomeBool {\IntCompare{3}>{2}}
\BoolVarLog \cFooSomeBool
```

\BoolSet *<boolean>* {*<boolexpr>*}

Evaluates the *<boolean expression>* and sets the *<boolean>* variable to the logical truth of this evaluation. For example

```
\BoolSet \lTmpaBool {\IntCompare{3}<{2}}
\BoolVarLog \lTmpaBool
```

\BoolSetTrue *<boolean>*

Sets *<boolean>* logically **true**.

\BoolSetFalse *<boolean>*

Sets *<boolean>* logically **false**.

\BoolSetEq *<boolean₁>* *<boolean₂>*

Sets *<boolean₁>* to the current value of *<boolean₂>*. For example

```
\BoolSetTrue \lTmpaBool
\BoolSetEq \lTmpbBool \lTmpaBool
\BoolVarLog \lTmpbBool
```

4.3 Viewing Booleans

\BoolLog {*<boolean expression>*}

Writes the logical truth of the *<boolean expression>* in the log file.

\BoolVarLog *<boolean>*

Writes the logical truth of the *<boolean>* in the log file.

\BoolShow {*<boolean expression>*}

Displays the logical truth of the *<boolean expression>* on the terminal.

\BoolVarShow *<boolean>*

Displays the logical truth of the *<boolean>* on the terminal.

4.4 Booleans and Conditionals

\BoolIfExist *<boolean>*

\BoolIfExistTF *<boolean>* {*<true code>*} {*<false code>*}

Tests whether the *<boolean>* is currently defined. This does not check that the *<boolean>* really is a boolean variable. For example

```
\BoolIfExistTF \lTmpaBool {\Result{Yes}} {\Result{No}}
\BoolIfExistTF \lFooUndefinedBool {\Result{Yes}} {\Result{No}}
```

Yes No

```
\BoolVarIf <boolean>
\BoolVarIfTF <boolean> {\<true code>} {\<false code>}
```

Tests the current truth of *<boolean>*, and continues evaluation based on this result. For example

```
\BoolSetTrue \lTmpaBool
\BoolVarIfTF \lTmpaBool {\Result{True!}} {\Result{False!}}
\BoolSetFalse \lTmpaBool
\BoolVarIfTF \lTmpaBool {\Result{True!}} {\Result{False!}}
```

True! False!

```
\BoolVarNot <boolean>
\BoolVarNotTF <boolean> {\<true code>} {\<false code>}
```

Evaluates *<true code>* if *<boolean>* is **false**, and *<false code>* If *<boolean>* is **true**. For example

```
\BoolVarNotTF {\IntCompare{3}>{2}} {\Result{Yes}} {\Result{No}}
```

No

```
\BoolVarAnd <boolean1> <boolean2>
\BoolVarAndTF <boolean1> <boolean2> {\<true code>} {\<false code>}
```

Implements the “And” operation between two booleans, hence is **true** if both are **true**. The *<boolean₂>* is only evaluated if it is needed to determine the result of **\BoolVarAnd**. For example

```
\BoolVarAndTF {\IntCompare{3}>{2}} {\IntCompare{3}>{4}} {\Result{Yes}} {\Result{No}}
```

No

```
\BoolVarOr <boolean1> <boolean2>
\BoolVarOrTF <boolean1> <boolean2> {\<true code>} {\<false code>}
```

Implements the “Or” operation between two booleans, hence is **true** if either one is **true**. The *<boolean₂>* is only evaluated if it is needed to determine the result of **\BoolVarOr**. For example

```
\BoolVarOrTF {\IntCompare{3}>{2}} {\IntCompare{3}>{4}} {\Result{Yes}} {\Result{No}}
```

Yes

```
\BoolVarXor <boolean1> <boolean2>
\BoolVarXorTF <boolean1> <boolean2> {\<true code>} {\<false code>}
```

Implements an “exclusive or” operation between two booleans. For example

```
\BoolVarXorTF {\IntCompare{3}>{2}} {\IntCompare{3}>{4}} {\Result{Yes}} {\Result{No}}
```

Yes

Chapter 5

Token Lists (Tl)

5.1 Constant and Scratch Token Lists

`\cSpaceTl`

An explicit space character contained in a token list. For use where an explicit space is required.

`\cEmptyTl`

Constant that is always empty.

`\lTmptaTl \lTmpbTl \lTmpcTl \lTmpiTl \lTmpjTl \lTmpkTl`

Scratch token lists for local assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

`\gTmptaTl \gTmpbTl \gTmpcTl \gTmpiTl \gTmpjTl \gTmpkTl`

Scratch token lists for global assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

5.2 Creating and Using Token Lists

`\TlNew` $\langle tl\ var \rangle$

Creates a new $\langle tl\ var \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle tl\ var \rangle$ is initially empty.

`\TlConst` $\langle tl\ var \rangle$ $\{ \langle token\ list \rangle \}$

Creates a new constant $\langle tl\ var \rangle$ or raises an error if the name is already taken. The value of the $\langle tl\ var \rangle$ is set globally to the $\langle token\ list \rangle$.

`\TlUse` $\langle tl\ var \rangle$

Recovers the content of a $\langle tl\ var \rangle$ and returns the value. An error is raised if the variable does not exist or if it is invalid. Note that it is possible to use a $\langle tl\ var \rangle$ directly without an accessor function.

 $\backslash\text{TlToStr}$ $\{\langle token\ list\rangle\}$

Converts the $\langle token\ list\rangle$ to a $\langle string\rangle$, returning the resulting character tokens. A $\langle string\rangle$ is a series of tokens with category code 12 (other) with the exception of spaces, which retain category code 10 (space).

 $\backslash\text{TlVarToStr}$ $\langle tl\ var\rangle$

Converts the content of the $\langle tl\ var\rangle$ to a string, returning the resulting character tokens. A $\langle string\rangle$ is a series of tokens with category code 12 (other) with the exception of spaces, which retain category code 10 (space).

5.3 Viewing Token Lists

 $\backslash\text{TlLog}$ $\{\langle token\ list\rangle\}$

Writes the $\langle token\ list\rangle$ in the log file. See also $\backslash\text{TlShow}$ which displays the result in the terminal.

 $\backslash\text{TlVarLog}$ $\langle tl\ var\rangle$

Writes the content of the $\langle tl\ var\rangle$ in the log file. See also $\backslash\text{TlVarShow}$ which displays the result in the terminal.

 $\backslash\text{TlShow}$ $\{\langle token\ list\rangle\}$

Displays the $\langle token\ list\rangle$ on the terminal.

 $\backslash\text{TlVarShow}$ $\langle tl\ var\rangle$

Displays the content of the $\langle tl\ var\rangle$ on the terminal. This is similar to the \TeX primitive $\backslash\text{show}$, wrapped to a fixed number of characters per line.

5.4 Setting Token List Variables

 $\backslash\text{TlSet}$ $\langle tl\ var\rangle\ \{\langle tokens\rangle\}$

Sets $\langle tl\ var\rangle$ to contain $\langle tokens\rangle$, removing any previous content from the variable. For example

```
\TlSet \lTmpiTl {\IntMathMult{4}{5}}
\TlUse \lTmpiTl
```

20

 $\backslash\text{TlSetEq}$ $\langle tl\ var_1\rangle\ \langle tl\ var_2\rangle$

Sets the content of $\langle tl\ var_1\rangle$ equal to that of $\langle tl\ var_2\rangle$.

 $\backslash\text{TlClear}$ $\langle tl\ var\rangle$

Clears all entries from the $\langle tl\ var\rangle$. For example


```
\TlSet \lTmptl {One}
\TlClear \lTmptl
\TlSet \lTmptl {Two}
\TlUse \lTmptl
```

Two

\TlClearNew $\langle tl\ var \rangle$

Ensures that the $\langle tl\ var \rangle$ exists globally by applying **\TlNew** if necessary, then applies **\TlClear** to leave the $\langle tl\ var \rangle$ empty.

\TlConcat $\langle tl\ var_1 \rangle \langle tl\ var_2 \rangle \langle tl\ var_3 \rangle$

Concatenates the content of $\langle tl\ var_2 \rangle$ and $\langle tl\ var_3 \rangle$ together and saves the result in $\langle tl\ var_1 \rangle$. The $\langle tl\ var_2 \rangle$ is placed at the left side of the new token list.

\TlPutLeft $\langle tl\ var \rangle \{ \langle tokens \rangle \}$

Appends $\langle tokens \rangle$ to the left side of the current content of $\langle tl\ var \rangle$. For example

```
\TlSet \lTmptl {Functional}
\TlPutLeft \lTmptl {Hello}
\TlUse \lTmptl
```

HelloFunctional

\TlPutRight $\langle tl\ var \rangle \{ \langle tokens \rangle \}$

Appends $\langle tokens \rangle$ to the right side of the current content of $\langle tl\ var \rangle$. For example

```
\TlSet \lTmptl {Functional}
\TlPutRight \lTmptl {World}
\TlUse \lTmptl
```

FunctionalWorld

5.5 Replacing Tokens

Within token lists, replacement takes place at the top level: there is no recursion into brace groups (more precisely, within a group defined by a category code 1/2 pair).

\TlReplaceOnce $\langle tl\ var \rangle \{ \langle old\ tokens \rangle \} \{ \langle new\ tokens \rangle \}$

Replaces the first (leftmost) occurrence of $\langle old\ tokens \rangle$ in the $\langle tl\ var \rangle$ with $\langle new\ tokens \rangle$. $\langle Old\ tokens \rangle$ cannot contain $\{$, $\}$ or $\#$ (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6).

\TlReplaceAll $\langle tl\ var \rangle \{ \langle old\ tokens \rangle \} \{ \langle new\ tokens \rangle \}$

Replaces all occurrences of $\langle old\ tokens \rangle$ in the $\langle tl\ var \rangle$ with $\langle new\ tokens \rangle$. $\langle Old\ tokens \rangle$ cannot contain $\{$, $\}$ or $\#$ (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6). As this function operates from left to right, the pattern $\langle old\ tokens \rangle$ may remain after the replacement (see **\TlRemoveAll** for an example).

\TlRemoveOnce $\langle tl\ var \rangle \{ \langle tokens \rangle \}$

Removes the first (leftmost) occurrence of $\langle tokens \rangle$ from the $\langle tl\ var \rangle$. $\langle Tokens \rangle$ cannot contain $\{$, $\}$ or $\#$ (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and

tokens with category code 6).

\TlRemoveAll $\langle tl\ var \rangle$ $\{\langle tokens \rangle\}$

Removes all occurrences of $\langle tokens \rangle$ from the $\langle tl\ var \rangle$. $\langle Tokens \rangle$ cannot contain $\{$, $\}$ or $\#$ (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6). As this function operates from left to right, the pattern $\langle tokens \rangle$ may remain after the removal, for instance,

```
\TlSet \lTmplTl {abbccd}
\TlRemoveAll \lTmplTl {bc}
\TlUse \lTmplTl
```

abcd

\TlTrimSpaces $\{\langle token\ list \rangle\}$

Removes any leading and trailing explicit space characters (explicit tokens with character code 32 and category code 10) from the $\langle token\ list \rangle$ and returns the result.

\TlVarTrimSpaces $\langle tl\ var \rangle$

Sets the $\langle tl\ var \rangle$ to contain the result of removing any leading and trailing explicit space characters (explicit tokens with character code 32 and category code 10) from its contents.

5.6 Working with the Content of Token Lists

\TlCount $\{\langle tokens \rangle\}$

Counts the number of $\langle items \rangle$ in $\langle tokens \rangle$ and returns this information. Unbraced tokens count as one element as do each token group $\{\dots\}$. This process ignores any unprotected spaces within $\langle tokens \rangle$.

\TlVarCount $\langle tl\ var \rangle$

Counts the number of $\langle items \rangle$ in the $\langle tl\ var \rangle$ and returns this information. Unbraced tokens count as one element as do each token group $\{\dots\}$. This process ignores any unprotected spaces within the $\langle tl\ var \rangle$.

\TlHead $\{\langle token\ list \rangle\}$

Returns the first $\langle item \rangle$ in the $\langle token\ list \rangle$, discarding the rest of the $\langle token\ list \rangle$. All leading explicit space characters (explicit tokens with character code 32 and category code 10) are discarded; for example

```
\fbox {1\TlHead{ abc }2}
\fbox {1\TlHead{ abc }2}
```

1a2 1a2

If the “head” is a brace group, rather than a single token, the braces are removed, and so

```
\TlHead { { ab} c }
```

yields $_ab$. A blank $\langle token\ list \rangle$ (see **\TlIfBlank**) results in **\TlHead** returning nothing.

\TlVarHead $\langle tl\ var \rangle$

Returns the first $\langle item \rangle$ in the $\langle tl\ var \rangle$, discarding the rest of the $\langle tl\ var \rangle$. All leading explicit space characters (explicit tokens with character code 32 and category code 10) are discarded.

\TlTail {*<token list>*}

Discards all leading explicit space characters (explicit tokens with character code 32 and category code 10) and the first *<item>* in the *<token list>*, and returns the remaining tokens. Thus for example

```
\TlTail { a {bc} d }
```

and

```
\TlTail { a {bc} d }
```

both return `{bc}{d}`. A blank *<token list>* (see **\TlIfBlank**) results in **\TlTail** returning nothing.

\TlVarTail *<tl var>*

Discards all leading explicit space characters (explicit tokens with character code 32 and category code 10) and the first *<item>* in the *<tl var>*, and returns the remaining tokens.

\TlItem {*<token list>*} {*<integer expression>*}

\TlVarItem *<tl var>* {*<integer expression>*}

Indexing items in the *<token list>* from 1 on the left, this function evaluates the *<integer expression>* and returns the appropriate item from the *<token list>*. If the *<integer expression>* is negative, indexing occurs from the right of the token list, starting at -1 for the right-most item. If the index is out of bounds, then the function returns nothing.

\TlRandItem {*<token list>*}

\TlVarRandItem *<tl var>*

Selects and returns a pseudo-random item of the *<token list>*. If the *<token list>* is blank, the result is empty.

5.7 Mapping over Token Lists

All mappings are done at the current group level, *i.e.* any local assignments made by the *<function>* or *<code>* discussed below remain in effect after the loop.

\TlMapInline {*<token list>*} {*<inline function>*}

Applies the *<inline function>* to every *<item>* stored within the *<token list>*. The *<inline function>* should consist of code which receives the *<item>* as **#1**.

\TlVarMapInline *<tl var>* {*<inline function>*}

Applies the *<inline function>* to every *<item>* stored within the *<tl var>*. The *<inline function>* should consist of code which receives the *<item>* as **#1**.

\TlMapVariable {*<token list>*} *<variable>* {*<code>*}

Stores each *<item>* of the *<token list>* in turn in the (token list) *<variable>* and applies the *<code>*. The *<code>* will usually make use of the *<variable>*, but this is not enforced. The assignments to the *<variable>* are local. Its value after the loop is the last *<item>* in the *<tl var>*, or its original value if the *<tl var>* is blank.

\TlVarMapVariable $\langle tl\ var \rangle$ $\langle variable \rangle$ $\{\langle code \rangle\}$

Stores each $\langle item \rangle$ of the $\langle tl\ var \rangle$ in turn in the (token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle item \rangle$ in the $\langle tl\ var \rangle$, or its original value if the $\langle tl\ var \rangle$ is blank.

5.8 Token List Conditionals

\TlIfExist $\langle tl\ var \rangle$
\TlIfExistTF $\langle tl\ var \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests whether the $\langle tl\ var \rangle$ is currently defined. This does not check that the $\langle tl\ var \rangle$ really is a token list variable.

\TlIfEmpty $\{\langle token\ list \rangle\}$
\TlIfEmptyTF $\{\langle token\ list \rangle\}$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests if the $\langle token\ list \rangle$ is entirely empty (*i.e.* contains no tokens at all). For example

```
\TlIfEmptyTF {abc} {\Result{Empty}} {\Result{NonEmpty}}
\TlIfEmptyTF {} {\Result{Empty}} {\Result{NonEmpty}}
```

NonEmpty Empty

\TlVarIfEmpty $\langle tl\ var \rangle$
\TlVarIfEmptyTF $\langle tl\ var \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests if the $\langle token\ list\ variable \rangle$ is entirely empty (*i.e.* contains no tokens at all). For example

```
\TlSet \lTmplTl {abc}
\TlVarIfEmptyTF \lTmplTl {\Result{Empty}} {\Result{NonEmpty}}
\TlClear \lTmplTl
\TlVarIfEmptyTF \lTmplTl {\Result{Empty}} {\Result{NonEmpty}}
```

NonEmpty Empty

\TlIfBlank $\{\langle token\ list \rangle\}$
\TlIfBlankTF $\{\langle token\ list \rangle\}$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests if the $\langle token\ list \rangle$ consists only of blank spaces (*i.e.* contains no item). The test is **true** if $\langle token\ list \rangle$ is zero or more explicit space characters (explicit tokens with character code 32 and category code 10), and is **false** otherwise.

\TlIfEq $\{\langle token\ list_1 \rangle\}$ $\{\langle token\ list_2 \rangle\}$
\TlIfEqTF $\{\langle token\ list_1 \rangle\}$ $\{\langle token\ list_2 \rangle\}$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Tests if $\langle token\ list_1 \rangle$ and $\langle token\ list_2 \rangle$ contain the same list of tokens, both in respect of character codes and category codes. See **\StrIfEq** if category codes are not important. For example

```
\TlIfEqTF {abc} {abc} {\Result{Yes}} {\Result{No}}
\TlIfEqTF {abc} {xyz} {\Result{Yes}} {\Result{No}}
```

Yes No

\TlVarIfEq $\langle tl\ var_1 \rangle$ $\langle tl\ var_2 \rangle$
\TlVarIfEqTF $\langle tl\ var_1 \rangle$ $\langle tl\ var_2 \rangle$ $\{\langle true\ code \rangle\}$ $\{\langle false\ code \rangle\}$

Compares the content of two $\langle token\ list\ variables \rangle$ and is logically **true** if the two contain the same

list of tokens (*i.e.* identical in both the list of characters they contain and the category codes of those characters). For example

```
\TlSet \lTmptaTl {abc}
\TlSet \lTmpbTl {abc}
\TlSet \lTmpcTl {xyz}
\TlVarIfEqTF \lTmptaTl \lTmpbTl {\Result{Yes}} {\Result{No}}
\TlVarIfEqTF \lTmptaTl \lTmpcTl {\Result{Yes}} {\Result{No}}
```

Yes No

See also `\StrVarIfEq` for a comparison that ignores category codes.

```
\TlIfIn {<token list1>} {<token list2>}
\TlIfInTF {<token list1>} {<token list2>} {<true code>} {<false code>}
```

Tests if $\langle token list_2 \rangle$ is found inside $\langle token list_1 \rangle$. The $\langle token list_2 \rangle$ cannot contain the tokens `{`, `}` or `#` (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6). The search does not enter brace (category code 1/2) groups.

```
\TlVarIfIn <tl var> {<token list>}
\TlVarIfInTF <tl var> {<token list>} {<true code>} {<false code>}
```

Tests if the $\langle token list \rangle$ is found in the content of the $\langle tl var \rangle$. The $\langle token list \rangle$ cannot contain the tokens `{`, `}` or `#` (more precisely, explicit character tokens with category code 1 (begin-group) or 2 (end-group), and tokens with category code 6).

```
\TlIfSingle {<token list>}
\TlIfSingleTF {<token list>} {<true code>} {<false code>}
```

Tests if the $\langle token list \rangle$ has exactly one $\langle item \rangle$, *i.e.* is a single normal token (neither an explicit space character nor a begin-group character) or a single brace group, surrounded by optional spaces on both sides. In other words, such a token list has token count 1 according to `\TlCount`.

```
\TlVarIfSingle <tl var>
\TlVarIfSingleTF <tl var> {<true code>} {<false code>}
```

Tests if the content of the $\langle tl var \rangle$ consists of a single $\langle item \rangle$, *i.e.* is a single normal token (neither an explicit space character nor a begin-group character) or a single brace group, surrounded by optional spaces on both sides. In other words, such a token list has token count 1 according to `\TlVarCount`.

5.9 Token List Case Functions

```
\TlVarCase <test token list variable>
{
  <token list variable case1> {<code case1>}
  <token list variable case2> {<code case2>}
  ...
  <token list variable casen> {<code casen>}
}
```

This function compares the $\langle test token list variable \rangle$ in turn with each of the $\langle token list variable cases \rangle$. If the two are equal (as described for `\TlVarIfEq`) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. The function does nothing if there is no match.

```

\TlVarCaseT <test token list variable>
{
  <token list variable case1> {<code case1>}
  <token list variable case2> {<code case2>}
  ...
  <token list variable casen> {<code casen>}
}
{<true code>}

```

This function compares the *<test token list variable>* in turn with each of the *<token list variable cases>*. If the two are equal (as described for **\TlVarIfEq**) then the associated *<code>* is left in the input stream and other cases are discarded. If any of the cases are matched, the *<true code>* is also inserted into the input stream (after the code for the appropriate case).

```

\TlVarCaseF <test token list variable>
{
  <token list variable case1> {<code case1>}
  <token list variable case2> {<code case2>}
  ...
  <token list variable casen> {<code casen>}
}
{<false code>}

```

This function compares the *<test token list variable>* in turn with each of the *<token list variable cases>*. If the two are equal (as described for **\TlVarIfEq**) then the associated *<code>* is left in the input stream and other cases are discarded. If none match then the *<false code>* is inserted into the input stream (after the code for the appropriate case).

```

\TlVarCaseTF <test token list variable>
{
  <token list variable case1> {<code case1>}
  <token list variable case2> {<code case2>}
  ...
  <token list variable casen> {<code casen>}
}
{<true code>}
{<false code>}

```

This function compares the *<test token list variable>* in turn with each of the *<token list variable cases>*. If the two are equal (as described for **\TlVarIfEq**) then the associated *<code>* is left in the input stream and other cases are discarded. If any of the cases are matched, the *<true code>* is also inserted into the input stream (after the code for the appropriate case), while if none match then the *<false code>* is inserted. The function **\TlVarCase**, which does nothing if there is no match, is also available.

Chapter 6

Strings (Str)

6.1 Constant and Scratch Strings

```
\cAmpersandStr \cAtsignStr \cBackslashStr \cLeftBraceStr \cRightBraceStr  
\cCircumflexStr \cColonStr \cDollarStr \cHashStr \cPercentStr \cTildeStr  
\cUnderscoreStr \cZeroStr
```

Constant strings, containing a single character token, with category code 12.

```
\lTmPaStr \lTmPbStr \lTmPcStr \lTmPiStr \lTmPjStr \lTmPkStr
```

Scratch strings for local assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

```
\gTmPaStr \gTmPbStr \gTmPcStr \gTmPiStr \gTmPjStr \gTmPkStr
```

Scratch strings for global assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

6.2 Creating and Using Strings

```
\StrNew <str var>
```

Creates a new <str var> or raises an error if the name is already taken. The declaration is global. The <str var> is initially empty.

```
\StrConst <str var> {<token list>}
```

Creates a new constant <str var> or raises an error if the name is already taken. The value of the <str var> is set globally to the <token list>, converted to a string.

```
\StrUse <str var>
```

Recovers the content of a <str var> and returns the value. An error is raised if the variable does not exist or if it is invalid. Note that it is possible to use a <str> directly without an accessor function.

6.3 Viewing Strings

\StrLog {*token list*}

Writes *token list* in the log file.

\StrVarLog *str var*

Writes the content of the *str var* in the log file. For example

```
\StrSet \lTmpiStr {1234\abcd5678}
\StrVarLog \lTmpiStr
```

\StrShow {*token list*}

Displays *token list* on the terminal.

\StrVarShow *str var*

Displays the content of the *str var* on the terminal.

6.4 Setting String Variables

\StrSet *str var* {*token list*}

Converts the *token list* to a *string*, and stores the result in *str var*. For example

```
\StrSet \lTmpiStr {\IntMathMult{4}{5}}
\StrUse \lTmpiStr
```

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\StrSetEq *str var*₁ *str var*₂

Sets the content of *str var*₁ equal to that of *str var*₂.

\StrClear *str var*

Clears the content of the *str var*. For example

```
\StrSet \lTnpjStr {One}
\StrClear \lTnpjStr
\StrSet \lTnpjStr {Two}
\StrUse \lTnpjStr
```

Two

\StrClearNew *str var*

Ensures that the *str var* exists globally by applying **\StrNew** if necessary, then applies **\StrClear** to leave the *str var* empty.

\StrConcat $\langle str\ var_1 \rangle\ \langle str\ var_2 \rangle\ \langle str\ var_3 \rangle$

Concatenates the content of $\langle str\ var_2 \rangle$ and $\langle str\ var_3 \rangle$ together and saves the result in $\langle str\ var_1 \rangle$. The $\langle str\ var_2 \rangle$ is placed at the left side of the new string variable. The $\langle str\ var_2 \rangle$ and $\langle str\ var_3 \rangle$ must indeed be strings, as this function does not convert their contents to a string.

\StrPutLeft $\langle str\ var \rangle\ \{\langle token\ list \rangle\}$

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$, and prepends the result to $\langle str\ var \rangle$. The current contents of the $\langle str\ var \rangle$ are not automatically converted to a string. For example

```
\StrSet \lTmkpStr {Functional}
\StrPutLeft \lTmkpStr {Hello}
\StrUse \lTmkpStr
```

HelloFunctional

\StrPutRight $\langle str\ var \rangle\ \{\langle token\ list \rangle\}$

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$, and appends the result to $\langle str\ var \rangle$. The current contents of the $\langle str\ var \rangle$ are not automatically converted to a string. For example

```
\StrSet \lTmkpStr {Functional}
\StrPutRight \lTmkpStr {World}
\StrUse \lTmkpStr
```

FunctionalWorld

6.5 Modifying String Variables

\StrReplaceOnce $\langle str\ var \rangle\ \{\langle old \rangle\}\ \{\langle new \rangle\}$

Converts the $\langle old \rangle$ and $\langle new \rangle$ token lists to strings, then replaces the first (leftmost) occurrence of $\langle old\ string \rangle$ in the $\langle str\ var \rangle$ with $\langle new\ string \rangle$.

\StrReplaceAll $\langle str\ var \rangle\ \{\langle old \rangle\}\ \{\langle new \rangle\}$

Converts the $\langle old \rangle$ and $\langle new \rangle$ token lists to strings, then replaces all occurrences of $\langle old\ string \rangle$ in the $\langle str\ var \rangle$ with $\langle new\ string \rangle$. As this function operates from left to right, the pattern $\langle old\ string \rangle$ may remain after the replacement.

\StrRemoveOnce $\langle str\ var \rangle\ \{\langle token\ list \rangle\}$

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$ then removes the first (leftmost) occurrence of $\langle string \rangle$ from the $\langle str\ var \rangle$.

\StrRemoveAll $\langle str\ var \rangle\ \{\langle token\ list \rangle\}$

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$ then removes all occurrences of $\langle string \rangle$ from the $\langle str\ var \rangle$. As this function operates from left to right, the pattern $\langle string \rangle$ may remain after the removal, for instance,

```
\StrSet \lTmpaStr {abbccd}
\StrRemoveAll \lTmpaStr {bc}
\TlUse \lTmpaStr
```

abcd

6.6 Working with the Content of Strings

\StrCount { $\langle token list \rangle$ }

Returns the number of characters in the string representation of $\langle token list \rangle$, as an integer denotation. All characters including spaces are counted.

Due to naming conflict, you need to use **\StrSize** instead of **\StrCount** if you want to use **functional** package together with **xstring** package.

\StrVarCount $\langle tl var \rangle$

Returns the number of characters in the string representation of the $\langle tl var \rangle$, as an integer denotation. All characters including spaces are counted.

\StrHead { $\langle token list \rangle$ }

Converts the $\langle token list \rangle$ into a $\langle string \rangle$. The first character in the $\langle string \rangle$ is then returned, with category code “other”. If the first character is a space, it returns a space token with category code 10 (blank space). If the $\langle string \rangle$ is empty, then nothing is returned.

\StrVarHead $\langle tl var \rangle$

Converts the $\langle tl var \rangle$ into a $\langle string \rangle$. The first character in the $\langle string \rangle$ is then returned, with category code “other”. If the first character is a space, it returns a space token with category code 10 (blank space). If the $\langle string \rangle$ is empty, then nothing is returned.

\StrTail { $\langle token list \rangle$ }

Converts the $\langle token list \rangle$ to a $\langle string \rangle$, removes the first character, and returns the remaining characters (if any) with category codes 12 and 10 (for spaces). If the first character is a space, it only trims that space. If the $\langle token list \rangle$ is empty, then nothing is left on the input stream.

\StrVarTail $\langle tl var \rangle$

Converts the $\langle tl var \rangle$ to a $\langle string \rangle$, removes the first character, and returns the remaining characters (if any) with category codes 12 and 10 (for spaces). If the first character is a space, it only trims that space. If the $\langle token list \rangle$ is empty, then nothing is left on the input stream.

\StrItem { $\langle token list \rangle$ } { $\langle integer expression \rangle$ }

Converts the $\langle token list \rangle$ to a $\langle string \rangle$, and returns the character in position $\langle integer expression \rangle$ of the $\langle string \rangle$, starting at 1 for the first (left-most) character. All characters including spaces are taken into account. If the $\langle integer expression \rangle$ is negative, characters are counted from the end of the $\langle string \rangle$. Hence, -1 is the right-most character, *etc.*

\StrVarItem $\langle tl var \rangle$ { $\langle integer expression \rangle$ }

Converts the $\langle tl var \rangle$ to a $\langle string \rangle$, and returns the character in position $\langle integer expression \rangle$ of the $\langle string \rangle$, starting at 1 for the first (left-most) character. All characters including spaces are taken into account. If the $\langle integer expression \rangle$ is negative, characters are counted from the end of the $\langle string \rangle$. Hence, -1 is the right-most character, *etc.*

6.7 Mapping over Strings

All mappings are done at the current group level, *i.e.* any local assignments made by the $\langle function \rangle$ or $\langle code \rangle$ discussed below remain in effect after the loop.

```
\StrMapInline { $\langle token list \rangle$ } { $\langle inline function \rangle$ }
\StrVarMapInline  $\langle str var \rangle$  { $\langle inline function \rangle$ }
```

Converts the $\langle token list \rangle$ to a $\langle string \rangle$ then applies the $\langle inline function \rangle$ to every $\langle character \rangle$ in the $\langle str var \rangle$ including spaces. The $\langle inline function \rangle$ should consist of code which receives the $\langle character \rangle$ as #1.

```
\StrMapVariable { $\langle token list \rangle$ }  $\langle variable \rangle$  { $\langle code \rangle$ }
\StrVarMapVariable  $\langle str var \rangle$   $\langle variable \rangle$  { $\langle code \rangle$ }
```

Converts the $\langle token list \rangle$ to a $\langle string \rangle$ then stores each $\langle character \rangle$ in the $\langle string \rangle$ (including spaces) in turn in the (string or token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle character \rangle$ in the $\langle string \rangle$, or its original value if the $\langle string \rangle$ is empty.

6.8 String Conditionals

```
\StrIfExist  $\langle str var \rangle$ 
\StrIfExistTF  $\langle str var \rangle$  { $\langle true code \rangle$ } { $\langle false code \rangle$ }
```

Tests whether the $\langle str var \rangle$ is currently defined. This does not check that the $\langle str var \rangle$ really is a string.

```
\StrVarIfEmpty  $\langle str var \rangle$ 
\StrVarIfEmptyTF  $\langle str var \rangle$  { $\langle true code \rangle$ } { $\langle false code \rangle$ }
```

Tests if the $\langle string variable \rangle$ is entirely empty (*i.e.* contains no characters at all). For example

```
\StrSet \lTmPaStr {abc}
\StrVarIfEmptyTF \lTmPaStr {\Result{Empty}} {\Result{NonEmpty}}
\StrClear \lTmPaStr
\StrVarIfEmptyTF \lTmPaStr {\Result{Empty}} {\Result{NonEmpty}}
```

NonEmpty Empty

```
\StrIfEq { $\langle tl_1 \rangle$ } { $\langle tl_2 \rangle$ }
\StrIfEqTF { $\langle tl_1 \rangle$ } { $\langle tl_2 \rangle$ } { $\langle true code \rangle$ } { $\langle false code \rangle$ }
```

Compares the two $\langle token lists \rangle$ on a character by character basis (namely after converting them to strings), and is **true** if the two $\langle strings \rangle$ contain the same characters in the same order. See **\TlIfEq** to compare tokens (including their category codes) rather than characters. For example

```
\StrIfEqTF {abc} {abc} {\Result{Yes}} {\Result{No}}
\StrIfEqTF {abc} {xyz} {\Result{Yes}} {\Result{No}}
```

Yes No

```
\StrVarIfEq  $\langle str var_1 \rangle$   $\langle str var_2 \rangle$ 
\StrVarIfEqTF  $\langle str var_1 \rangle$   $\langle str var_2 \rangle$  { $\langle true code \rangle$ } { $\langle false code \rangle$ }
```

Compares the content of two $\langle str variables \rangle$ and is logically **true** if the two contain the same characters in the same order. See **\TlVarIfEq** to compare tokens (including their category codes) rather than characters.

```

\StrSet \lTmPaStr {abc}
\StrSet \lTmPbStr {abc}
\StrSet \lTmPcStr {xyz}
\StrVarIfEqTF \lTmPaStr \lTmPbStr {\Result{Yes}} {\Result{No}}
\StrVarIfEqTF \lTmPaStr \lTmPcStr {\Result{Yes}} {\Result{No}}

```

Yes	No
-----	----

```

\StrIfInTF {\<tl1>} {\<tl2>}
\StrIfInTF {\<tl1>} {\<tl2>} {\<true code>} {\<false code>}

```

Converts both $\langle token\ lists \rangle$ to $\langle strings \rangle$ and tests whether $\langle string_2 \rangle$ is found inside $\langle string_1 \rangle$.

```

\StrVarIfInTF \<str var> {\<token list>}
\StrVarIfInTF \<str var> {\<token list>} {\<true code>} {\<false code>}

```

Converts the $\langle token\ list \rangle$ to a $\langle string \rangle$ and tests if that $\langle string \rangle$ is found in the content of the $\langle str\ var \rangle$.

```

\StrCompare {\<tl1>} \<relation> {\<tl2>}
\StrCompareTF {\<tl1>} \<relation> {\<tl2>} {\<true code>} {\<false code>}

```

Compares the two $\langle token\ lists \rangle$ on a character by character basis (namely after converting them to strings) in a lexicographic order according to the character codes of the characters. The $\langle relation \rangle$ can be $<$, $=$, or $>$ and the test is **true** under the following conditions:

- for $<$, if the first string is earlier than the second in lexicographic order;
- for $=$, if the two strings have exactly the same characters;
- for $>$, if the first string is later than the second in lexicographic order.

For example:

```

\StrCompareTF {ab} < {abc} {\Result{Yes}} {\Result{No}}
\StrCompareTF {ab} < {aa} {\Result{Yes}} {\Result{No}}

```

Yes	No
-----	----

Due to naming conflict, you need to use `\StrIfCompare`/`\StrIfCompareTF` as a replacement if you want to use **functional** package together with **xstring** package.

6.9 String Case Functions

```

\StrCase {\<test string>}
{
  {\<string case1>} {\<code case1>}
  {\<string case2>} {\<code case2>}
  ...
  {\<string casen>} {\<code casen>}
}

```

Compares the $\langle test\ string \rangle$ in turn with each of the $\langle string\ cases \rangle$ (all token lists are converted to strings). If the two are equal (as described for `\StrIfEq`) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded.

```

\StrCaseT {⟨test string⟩}
{
  {⟨string case1⟩} {⟨code case1⟩}
  {⟨string case2⟩} {⟨code case2⟩}
  ...
  {⟨string casen⟩} {⟨code casen⟩}
}
{⟨true code⟩}

```

Compares the $\langle test\ string \rangle$ in turn with each of the $\langle string\ cases \rangle$ (all token lists are converted to strings). If the two are equal (as described for $\backslash StrIfEq$) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case).

```

\StrCaseF {⟨test string⟩}
{
  {⟨string case1⟩} {⟨code case1⟩}
  {⟨string case2⟩} {⟨code case2⟩}
  ...
  {⟨string casen⟩} {⟨code casen⟩}
}
{⟨false code⟩}

```

Compares the $\langle test\ string \rangle$ in turn with each of the $\langle string\ cases \rangle$ (all token lists are converted to strings). If the two are equal (as described for $\backslash StrIfEq$) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If none match then the $\langle false\ code \rangle$ is inserted.

```

\StrCaseTF {⟨test string⟩}
{
  {⟨string case1⟩} {⟨code case1⟩}
  {⟨string case2⟩} {⟨code case2⟩}
  ...
  {⟨string casen⟩} {⟨code casen⟩}
}
{⟨true code⟩}
{⟨false code⟩}

```

Compares the $\langle test\ string \rangle$ in turn with each of the $\langle string\ cases \rangle$ (all token lists are converted to strings). If the two are equal (as described for $\backslash StrIfEq$) then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted.

Chapter 7

Integers (Int)

7.1 Constant and Scratch Integers

`\cZeroInt \cOneInt`

Integer values used with primitive tests and assignments: their self-terminating nature makes these more convenient and faster than literal numbers.

`\cMaxInt`

The maximum value that can be stored as an integer.

`\cMaxRegisterInt`

Maximum number of registers.

`\cMaxCharInt`

Maximum character code completely supported by the engine.

`\lTmPaInt \lTmPbInt \lTmPcInt \lTmPiInt \lTmPjInt \lTmPkInt`

Scratch integer for local assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

`\gTmPaInt \gTmPbInt \gTmPcInt \gTmPiInt \gTmPjInt \gTmPkInt`

Scratch integer for global assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

7.2 Integer Expressions

`\IntEval {⟨integer expression⟩}`

Evaluates the *⟨integer expression⟩* and returns the result: for positive results an explicit sequence of decimal digits not starting with 0, for negative results – followed by such a sequence, and 0 for zero. For example

```
\IntEval {(1+4)*(2-3)/5}
```

-1

```
\IntMathAdd {⟨integer expression1⟩} {⟨integer expression2⟩}
```

Adds {⟨integer expression₁⟩} and {⟨integer expression₂⟩}, and returns the result. For example

```
\IntMathAdd {7} {3}
```

10

```
\IntMathSub {⟨integer expression1⟩} {⟨integer expression2⟩}
```

Subtracts {⟨integer expression₂⟩} from {⟨integer expression₁⟩}, and returns the result. For example

```
\IntMathSub {7} {3}
```

4

```
\IntMathMult {⟨integer expression1⟩} {⟨integer expression2⟩}
```

Multiplies {⟨integer expression₁⟩} by {⟨integer expression₂⟩}, and returns the result. For example

```
\IntMathMult {7} {3}
```

21

```
\IntMathDiv {⟨integer expression1⟩} {⟨integer expression2⟩}
```

Evaluates the two ⟨integer expressions⟩ as described earlier, then divides the first value by the second, and rounds the result to the closest integer. Ties are rounded away from zero. Note that this is identical to using / directly in an ⟨integer expression⟩. The result is returned as an ⟨integer denotation⟩. For example

```
\IntMathDiv {8} {3}
```

3

```
\IntMathDivTruncate {⟨integer expression1⟩} {⟨integer expression2⟩}
```

Evaluates the two ⟨integer expressions⟩ as described earlier, then divides the first value by the second, and rounds the result towards zero. Note that division using / rounds to the closest integer instead. The result is returned as an ⟨integer denotation⟩. For example

```
\IntMathDivTruncate {8} {3}
```

2

```
\IntMathSign {⟨intexpr⟩}
```

Evaluates the ⟨integer expression⟩ then leaves 1 or 0 or −1 in the input stream according to the sign of the result.

```
\IntMathAbs {⟨integer expression⟩}
```

Evaluates the ⟨integer expression⟩ as described for **\IntEval** and leaves the absolute value of the result in the input stream as an ⟨integer denotation⟩ after two expansions.

```
\IntMathMax {⟨intexpr1⟩} {⟨intexpr2⟩}
```

```
\IntMathMin {⟨intexpr1⟩} {⟨intexpr2⟩}
```

Evaluates the ⟨integer expressions⟩ as described for **\IntEval** and leaves either the larger or smaller value in the input stream as an ⟨integer denotation⟩ after two expansions.

 $\backslash\text{IntMathMod } \{\langle integer \rangle\} \{\langle integer \rangle\}$

Evaluates the two $\langle integer \rangle$ s as described earlier, then calculates the integer remainder of dividing the first expression by the second. This is obtained by subtracting $\backslash\text{IntMathDivTruncate } \{\langle integer \rangle\} \{\langle integer \rangle\}$ times $\langle integer \rangle$ from $\langle integer \rangle$. Thus, the result has the same sign as $\langle integer \rangle$ and its absolute value is strictly less than that of $\langle integer \rangle$. The result is left in the input stream as an $\langle integer \rangle$ after two expansions.

 $\backslash\text{IntMathRand } \{\langle integer \rangle\} \{\langle integer \rangle\}$

Evaluates the two $\langle integer \rangle$ s and produces a pseudo-random number between the two (with bounds included).

7.3 Creating and Using Integers

 $\backslash\text{IntNew } \langle integer \rangle$

Creates a new $\langle integer \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle integer \rangle$ is initially equal to 0.

 $\backslash\text{IntConst } \langle integer \rangle \{\langle integer \rangle\}$

Creates a new constant $\langle integer \rangle$ or raises an error if the name is already taken. The value of the $\langle integer \rangle$ is set globally to the $\langle integer \rangle$.

 $\backslash\text{IntUse } \langle integer \rangle$

Recovers the content of an $\langle integer \rangle$ and returns the value. An error is raised if the variable does not exist or if it is invalid.

7.4 Viewing Integers

 $\backslash\text{IntLog } \{\langle integer \rangle\}$

Writes the result of evaluating the $\langle integer \rangle$ in the log file.

 $\backslash\text{IntVarLog } \langle integer \rangle$

Writes the value of the $\langle integer \rangle$ in the log file.

 $\backslash\text{IntShow } \{\langle integer \rangle\}$

Displays the result of evaluating the $\langle integer \rangle$ on the terminal.

 $\backslash\text{IntVarShow } \langle integer \rangle$

Displays the value of the $\langle integer \rangle$ on the terminal.

7.5 Setting Integer Variables

\IntSet $\langle integer \rangle$ { $\langle integer expression \rangle$ }

Sets $\langle integer \rangle$ to the value of $\langle integer expression \rangle$, which must evaluate to an integer (as described for **\IntEval**). For example

```
\IntSet \lTmPaInt {3+5}
\IntUse \lTmPaInt
```

8

\IntSetEq $\langle integer_1 \rangle$ $\langle integer_2 \rangle$

Sets the content of $\langle integer_1 \rangle$ equal to that of $\langle integer_2 \rangle$.

\IntZero $\langle integer \rangle$

Sets $\langle integer \rangle$ to 0. For example

```
\IntSet \lTmPaInt {5}
\IntZero \lTmPaInt
\IntUse \lTmPaInt
```

0

\IntZeroNew $\langle integer \rangle$

Ensures that the $\langle integer \rangle$ exists globally by applying **\IntNew** if necessary, then applies **\IntZero** to leave the $\langle integer \rangle$ set to zero.

\IntIncr $\langle integer \rangle$

Increases the value stored in $\langle integer \rangle$ by 1. For example

```
\IntSet \lTmPaInt {5}
\IntIncr \lTmPaInt
\IntUse \lTmPaInt
```

6

\IntDecr $\langle integer \rangle$

Decreases the value stored in $\langle integer \rangle$ by 1. For example

```
\IntSet \lTmPaInt {5}
\IntDecr \lTmPaInt
\IntUse \lTmPaInt
```

4

\IntAdd $\langle integer \rangle$ { $\langle integer expression \rangle$ }

Adds the result of the $\langle integer expression \rangle$ to the current content of the $\langle integer \rangle$. For example

```
\IntSet \lTmPaInt {5}
\IntAdd \lTmPaInt {2}
\IntUse \lTmPaInt
```

7

\IntSub $\langle integer \rangle$ $\{\langle integer expression \rangle\}$

Subtracts the result of the $\langle integer expression \rangle$ from the current content of the $\langle integer \rangle$. For example

```
\IntSet \lTmPaInt {5}
\IntSub \lTmPaInt {3}
\IntUse \lTmPaInt
```

2

7.6 Integer Step Functions

\IntStepInline $\{\langle initial value \rangle\} \{\langle step \rangle\} \{\langle final value \rangle\} \{\langle code \rangle\}$

This function first evaluates the $\langle initial value \rangle$, $\langle step \rangle$ and $\langle final value \rangle$, all of which should be integer expressions. Then for each $\langle value \rangle$ from the $\langle initial value \rangle$ to the $\langle final value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream with $\#1$ replaced by the current $\langle value \rangle$. Thus the $\langle code \rangle$ should define a function of one argument ($\#1$). For example

```
\IgnoreSpacesOn
\TlClear \lTmPaTl
\IntStepInline {1} {3} {30} {
  \TlPutRight \lTmPaTl {[#1]}
}
\Result {\Value\lTmPaTl}
\IgnoreSpacesOff
```

produces [1][4][7][10][13][16][19][22][25][28].

\IntStepVariable $\{\langle initial value \rangle\} \{\langle step \rangle\} \{\langle final value \rangle\} \langle tl var \rangle \{\langle code \rangle\}$

This function first evaluates the $\langle initial value \rangle$, $\langle step \rangle$ and $\langle final value \rangle$, all of which should be integer expressions. Then for each $\langle value \rangle$ from the $\langle initial value \rangle$ to the $\langle final value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is evaluated, with the $\langle tl var \rangle$ defined as the current $\langle value \rangle$. Thus the $\langle code \rangle$ should make use of the $\langle tl var \rangle$.

7.7 Integer Conditionals

\IntIfExist $\langle integer \rangle$
\IntIfExistTF $\langle integer \rangle \{\langle true code \rangle\} \{\langle false code \rangle\}$

Tests whether the $\langle integer \rangle$ is currently defined. This does not check that the $\langle integer \rangle$ really is an integer variable.

\IntIfOdd $\{\langle integer expression \rangle\}$
\IntIfOddTF $\{\langle integer expression \rangle\} \{\langle true code \rangle\} \{\langle false code \rangle\}$

This function first evaluates the $\langle integer expression \rangle$ as described for **\IntEval**. It then evaluates if this is odd or even, as appropriate.

\IntIfEven $\{\langle integer expression \rangle\}$
\IntIfEvenTF $\{\langle integer expression \rangle\} \{\langle true code \rangle\} \{\langle false code \rangle\}$

This function first evaluates the $\langle integer expression \rangle$ as described for **\IntEval**. It then evaluates if this is even or odd, as appropriate.

\IntCompare {<integer expression>} <relation> {<integer expression>}

\IntCompareTF {<integer expression>} <relation> {<integer expression>} {<true code>} {<false code>}

This function first evaluates each of the <integer expressions> as described for **\IntEval**. The two results are then compared using the <relation>:

Equal	=
Greater than	>
Less than	<

For example

```
\IntCompareTF {2} > {1} {\Result{Greater}} {\Result{Less}}
\IntCompareTF {2} > {3} {\Result{Greater}} {\Result{Less}}
```

Greater Less

7.8 Integer Case Functions

\IntCase {<test integer expression>}

```
{
  {<integer case1>} {<code case1>}
  {<integer case2>} {<code case2>}
  ...
  {<integer casen>} {<code casen>}
}
```

This function evaluates the <test integer expression> and compares this in turn to each of the <integer expression cases>. If the two are equal then the associated <code> is left in the input stream and other cases are discarded.

\IntCaseT {<test integer expression>}

```
{
  {<integer case1>} {<code case1>}
  {<integer case2>} {<code case2>}
  ...
  {<integer casen>} {<code casen>}
}
{<true code>}
```

This function evaluates the <test integer expression> and compares this in turn to each of the <integer expression cases>. If the two are equal then the associated <code> is left in the input stream and other cases are discarded. If any of the cases are matched, the <true code> is also inserted into the input stream (after the code for the appropriate case).

\IntCaseF {<test integer expression>}

```
{
  {<integer case1>} {<code case1>}
  {<integer case2>} {<code case2>}
  ...
  {<integer casen>} {<code casen>}
}
{<false code>}
```

This function evaluates the <test integer expression> and compares this in turn to each of the <integer expression cases>. If the two are equal then the associated <code> is left in the input stream and other cases are discarded. If none match then the <false code> is into the input stream (after the code for the appropriate case). For example

```

\IgnoreSpacesOn
\IntCaseF { 2 * 5 }
{
  { 5 }      { Small }
  { 4 + 6 }  { Medium }
  { -2 * 10 } { Negative }
}
{ No idea! }
\IgnoreSpacesOff

```

Medium

```

\IntCaseTF {⟨test integer expression⟩}
{
  {⟨intexpr case1⟩} {⟨code case1⟩}
  {⟨intexpr case2⟩} {⟨code case2⟩}
  ...
  {⟨intexpr casen⟩} {⟨code casen⟩}
}
{⟨true code⟩}
{⟨false code⟩}

```

This function evaluates the *⟨test integer expression⟩* and compares this in turn to each of the *⟨integer expression cases⟩*. If the two are equal then the associated *⟨code⟩* is left in the input stream and other cases are discarded. If any of the cases are matched, the *⟨true code⟩* is also inserted into the input stream (after the code for the appropriate case), while if none match then the *⟨false code⟩* is inserted.

Chapter 8

Floating Point Numbers (Fp)

8.1 Constant and Scratch Floating Points

`\cZeroFp` `\cMinusZeroFp`

Zero, with either sign.

`\cOneFp`

One as an `fp`: useful for comparisons in some places.

`\cInfFp` `\cMinusInfFp`

Infinity, with either sign. These can be input directly in a floating point expression as `inf` and `-inf`.

`\cEFp`

The value of the base of the natural logarithm, $e = \exp(1)$.

`\cPiFp`

The value of π . This can be input directly in a floating point expression as `pi`.

`\cOneDegreeFp`

The value of 1° in radians. Multiply an angle given in degrees by this value to obtain a result in radians. Note that trigonometric functions expecting an argument in radians or in degrees are both available. Within floating point expressions, this can be accessed as `deg`.

`\lTmPaFp` `\lTmPbFp` `\lTmPcFp` `\lTmPiFp` `\lTmPjFp` `\lTmPkFp`

Scratch floating point numbers for local assignment. These are never used by the `functional` package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

`\gTmPaFp` `\gTmPbFp` `\gTmPcFp` `\gTmPiFp` `\gTmPjFp` `\gTmPkFp`

Scratch floating point numbers for global assignment. These are never used by the `functional` package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

8.2 Floating Point Expressions

\FpEval {<floating point expression>}

Evaluates the <floating point expression> and returns the result as a decimal number with no exponent. Leading or trailing zeros may be inserted to compensate for the exponent. Non-significant trailing zeros are trimmed, and integers are expressed without a decimal separator. The values $\pm\infty$ and NaN trigger an “invalid operation” exception. For a tuple, each item is converted using **\FpEval** and they are combined as $(\langle fp_1 \rangle, \sqcup \langle fp_2 \rangle, \sqcup \dots \langle fp_n \rangle)$ if $n > 1$ and $(\langle fp_1 \rangle,)$ or $()$ for fewer items. For example

```
\FpEval {(1.2+3.4)*(5.6-7.8)/9}
```

-1.124444444444444

\FpMathAdd {<fpexpr₁>} {<fpexpr₂>}

Adds {<fpexpr₁>} and {<fpexpr₂>}, and returns the result. For example

```
\FpMathAdd {2.8} {3.7}
\FpMathAdd {3.8-1} {2.7+1}
```

6.5 6.5

\FpMathSub {<fpexpr₁>} {<fpexpr₂>}

Subtracts {<fpexpr₂>} from {<fpexpr₁>}, and returns the result. For example

```
\FpMathSub {2.8} {3.7}
\FpMathSub {3.8-1} {2.7+1}
```

-0.9 -0.9

\FpMathMult {<fpexpr₁>} {<fpexpr₂>}

Multiplies {<fpexpr₁>} by {<fpexpr₂>}, and returns the result. For example

```
\FpMathMult {2.8} {3.7}
\FpMathMult {3.8-1} {2.7+1}
```

10.36 10.36

\FpMathDiv {<fpexpr₁>} {<fpexpr₂>}

Divides {<fpexpr₁>} by {<fpexpr₂>}, and returns the result. For example

```
\FpMathDiv {2.8} {3.7}
\FpMathDiv {3.8-1} {2.7+1}
```

0.7567567567567568 0.7567567567567568

\FpMathSign {<fpexpr>}

Evaluates the <fpexpr> and returns the value using **\FpEval{sign(<result>)}**: +1 for positive numbers and for $+\infty$, -1 for negative numbers and for $-\infty$, ± 0 for ± 0 . If the operand is a tuple or is NaN, then “invalid operation” occurs and the result is 0. For example

```
\FpMathSign {3.5}
\FpMathSign {-2.7}
```

1 -1

\FpMathAbs {*floating point expression*}

Evaluates the *floating point expression* as described for **\FpEval** and returns the absolute value. If the argument is $\pm\infty$, NaN or a tuple, “invalid operation” occurs. Within floating point expressions, **abs()** can be used; it accepts $\pm\infty$ and NaN as arguments.

\FpMathMax {*fp expression*₁} {*fp expression*₂}

\FpMathMin {*fp expression*₁} {*fp expression*₂}

Evaluates the *floating point expressions* as described for **\FpEval** and returns the resulting larger (**max**) or smaller (**min**) value. If the argument is a tuple, “invalid operation” occurs, but no other case raises exceptions. Within floating point expressions, **max()** and **min()** can be used.

8.3 Creating and Using Floating Points

\FpNew *fp var*

Creates a new *fp var* or raises an error if the name is already taken. The declaration is global. The *fp var* is initially +0.

\FpConst *fp var* {*floating point expression*}

Creates a new constant *fp var* or raises an error if the name is already taken. The *fp var* is set globally equal to the result of evaluating the *floating point expression*. For example

```
\FpConst \cMyPiFp {3.1415926}
\FpUse \cMyPiFp
```

3.1415926

\FpUse *fp var*

Recovers the value of the *fp var* and returns the value as a decimal number with no exponent.

8.4 Viewing Floating Points

\FpLog {*floating point expression*}

Evaluates the *floating point expression* and writes the result in the log file.

\FpVarLog *fp var*

Writes the value of *fp var* in the log file.

\FpShow {*floating point expression*}

Evaluates the *floating point expression* and displays the result in the terminal.

\FpVarShow *fp var*

Displays the value of *fp var* in the terminal.

8.5 Setting Floating Point Variables

\FpSet $\langle fp\ var \rangle \{ \langle floating\ point\ expression \rangle \}$

Sets $\langle fp\ var \rangle$ equal to the result of computing the $\langle floating\ point\ expression \rangle$. For example

```
\FpSet \lTmPaFp {4/7}
\FpUse \lTmPaFp
```

0.5714285714285714

\FpSetEq $\langle fp\ var_1 \rangle \langle fp\ var_2 \rangle$

Sets the floating point variable $\langle fp\ var_1 \rangle$ equal to the current value of $\langle fp\ var_2 \rangle$.

\FpZero $\langle fp\ var \rangle$

Sets the $\langle fp\ var \rangle$ to +0. For example

```
\FpSet \lTmPaFp {5.3}
\FpZero \lTmPaFp
\FpUse \lTmPaFp
```

0

\FpZeroNew $\langle fp\ var \rangle$

Ensures that the $\langle fp\ var \rangle$ exists globally by applying **\FpNew** if necessary, then applies **\FpZero** to leave the $\langle fp\ var \rangle$ set to +0.

\FpAdd $\langle fp\ var \rangle \{ \langle floating\ point\ expression \rangle \}$

Adds the result of computing the $\langle floating\ point\ expression \rangle$ to the $\langle fp\ var \rangle$. This also applies if $\langle fp\ var \rangle$ and $\langle floating\ point\ expression \rangle$ evaluate to tuples of the same size. For example

```
\FpSet \lTmPaFp {5.3}
\FpAdd \lTmPaFp {2.11}
\FpUse \lTmPaFp
```

7.41

\FpSub $\langle fp\ var \rangle \{ \langle floating\ point\ expression \rangle \}$

Subtracts the result of computing the $\langle floating\ point\ expression \rangle$ from the $\langle fp\ var \rangle$. This also applies if $\langle fp\ var \rangle$ and $\langle floating\ point\ expression \rangle$ evaluate to tuples of the same size. For example

```
\FpSet \lTmPaFp {5.3}
\FpSub \lTmPaFp {2.11}
\FpUse \lTmPaFp
```

3.19

8.6 Floating Point Step Functions

\FpStepInline $\{ \langle initial\ value \rangle \} \{ \langle step \rangle \} \{ \langle final\ value \rangle \} \{ \langle code \rangle \}$

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be floating point expressions evaluating to a floating point number, not a tuple. Then for each $\langle value \rangle$ from the

$\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream with **#1** replaced by the current $\langle value \rangle$. Thus the $\langle code \rangle$ should define a function of one argument (**#1**). For example

```
\IgnoreSpacesOn
\TlClear \lTmptl
\FpStepInline {1} {0.1} {1.5} {
  \TlPutRight \lTmptl {[#1]}
}
\Result {\Value\lTmptl}
\IgnoreSpacesOff
```

produces [1][1.1][1.2][1.3][1.4][1.5].

\FpStepVariable $\langle initial\ value \rangle$ $\langle step \rangle$ $\langle final\ value \rangle$ $\langle tl\ var \rangle$ $\langle code \rangle$

This function first evaluates the $\langle initial\ value \rangle$, $\langle step \rangle$ and $\langle final\ value \rangle$, all of which should be floating point expressions evaluating to a floating point number, not a tuple. Then for each $\langle value \rangle$ from the $\langle initial\ value \rangle$ to the $\langle final\ value \rangle$ in turn (using $\langle step \rangle$ between each $\langle value \rangle$), the $\langle code \rangle$ is inserted into the input stream, with the $\langle tl\ var \rangle$ defined as the current $\langle value \rangle$. Thus the $\langle code \rangle$ should make use of the $\langle tl\ var \rangle$.

8.7 Float Point Conditionals

\FpIfExist $\langle fp\ var \rangle$
\FpIfExistTF $\langle fp\ var \rangle$ $\langle true\ code \rangle$ $\langle false\ code \rangle$

Tests whether the $\langle fp\ var \rangle$ is currently defined. This does not check that the $\langle fp\ var \rangle$ really is a floating point variable. For example

```
\FpIfExistTF \lTmptl {\Result{Yes}} {\Result{No}}
\FpIfExistTF \lMyUndefinedFp {\Result{Yes}} {\Result{No}}
```

Yes No

\FpCompare $\langle fpexpr_1 \rangle$ $\langle relation \rangle$ $\langle fpexpr_2 \rangle$
\FpCompareTF $\langle fpexpr_1 \rangle$ $\langle relation \rangle$ $\langle fpexpr_2 \rangle$ $\langle true\ code \rangle$ $\langle false\ code \rangle$

Compares the $\langle fpexpr_1 \rangle$ and the $\langle fpexpr_2 \rangle$, and returns **true** if the $\langle relation \rangle$ is obeyed. For example

```
\FpCompareTF {1} > {0.9999} {\Result{Greater}} {\Result{Less}}
\FpCompareTF {1} > {1.0001} {\Result{Greater}} {\Result{Less}}
```

Greater Less

Two floating points x and y may obey four mutually exclusive relations: $x < y$, $x = y$, $x > y$, or $x?y$ (“not ordered”). The last case occurs exactly if one or both operands is **NaN** or is a tuple, unless they are equal tuples. Note that a **NaN** is distinct from any value, even another **NaN**, hence $x = x$ is not true for a **NaN**. To test if a value is **NaN**, compare it to an arbitrary number with the “not ordered” relation.

Tuples are equal if they have the same number of items and items compare equal (in particular there must be no **NaN**). At present any other comparison with tuples yields ? (not ordered). This is experimental.

Chapter 9

Dimensions (Dim)

9.1 Constant and Scratch Dimensions

`\cMaxDim`

The maximum value that can be stored as a dimension. This can also be used as a component of a skip.

`\cZeroDim`

A zero length as a dimension. This can also be used as a component of a skip.

`\lTmptaDim \lTmptbDim \lTmptcDim \lTmptiDim \lTmptjDim \lTmptkDim`

Scratch dimensions for local assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

`\gTmptaDim \gTmptbDim \gTmptcDim \gTmptiDim \gTmptjDim \gTmptkDim`

Scratch dimensions for global assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

9.2 Dimension Expressions

`\DimEval` $\{\langle dimension\ expression \rangle\}$

Evaluates the $\langle dimension\ expression \rangle$, expanding any dimensions and token list variables within the $\langle expression \rangle$ to their content (without requiring `\DimUse`/`\TlUse`) and applying the standard mathematical rules. The result of the calculation is returned as a $\langle dimension\ denotation \rangle$. For example

```
\DimEval {(1.2pt+3.4pt)/9}
```

0.51111pt

`\DimMathAdd` $\{\langle dimexpr_1 \rangle\} \{\langle dimexpr_2 \rangle\}$

Adds $\{\langle dimexpr_1 \rangle\}$ and $\{\langle dimexpr_2 \rangle\}$, and returns the result. For example

```
\DimMathAdd {2.8pt} {3.7pt}  
\DimMathAdd {3.8pt-1pt} {2.7pt+1pt}
```

6.5pt 6.5pt

 $\backslash\text{DimMathSub}$ $\{\langle\text{dimexpr}_1\rangle\}$ $\{\langle\text{dimexpr}_2\rangle\}$

Subtracts $\{\langle\text{dimexpr}_2\rangle\}$ from $\{\langle\text{dimexpr}_1\rangle\}$, and returns the result. For example

$\backslash\text{DimMathSub}$ $\{2.8\text{pt}\}$ $\{3.7\text{pt}\}$	-0.9pt
$\backslash\text{DimMathSub}$ $\{3.8\text{pt}-1\text{pt}\}$ $\{2.7\text{pt}+1\text{pt}\}$	-0.9pt

 $\backslash\text{DimMathRatio}$ $\{\langle\text{dimexpr}_1\rangle\}$ $\{\langle\text{dimexpr}_2\rangle\}$

Parses the two $\langle\text{dimension expressions}\rangle$, then calculates the ratio of the two and returns it. The result is a ratio expression between two integers, with all distances converted to scaled points. For example

$\backslash\text{DimMathRatio}$ $\{5\text{pt}\}$ $\{10\text{pt}\}$	$327680/655360$
--	-----------------

The returned value is suitable for use inside a $\langle\text{dimension expression}\rangle$ such as

$\backslash\text{DimSet}$ $\backslash\text{TmpaDim}$ $\{10\text{pt}*\backslash\text{DimMathRatio}\{5\text{pt}\}\{10\text{pt}\}\}$

 $\backslash\text{DimMathSign}$ $\{\langle\text{dimexpr}\rangle\}$

Evaluates the $\langle\text{dimexpr}\rangle$ then returns 1 or 0 or -1 according to the sign of the result. For example

$\backslash\text{DimMathSign}$ $\{3.5\text{pt}\}$	1
$\backslash\text{DimMathSign}$ $\{-2.7\text{pt}\}$	-1

 $\backslash\text{DimMathAbs}$ $\{\langle\text{dimexpr}\rangle\}$

Converts the $\langle\text{dimexpr}\rangle$ to its absolute value, returning the result as a $\langle\text{dimension denotation}\rangle$. For example

$\backslash\text{DimMathAbs}$ $\{3.5\text{pt}\}$	3.5pt
$\backslash\text{DimMathAbs}$ $\{-2.7\text{pt}\}$	2.7pt

 $\backslash\text{DimMathMax}$ $\{\langle\text{dimexpr}_1\rangle\}$ $\{\langle\text{dimexpr}_2\rangle\}$

 $\backslash\text{DimMathMin}$ $\{\langle\text{dimexpr}_1\rangle\}$ $\{\langle\text{dimexpr}_2\rangle\}$

Evaluates the two $\langle\text{dimension expressions}\rangle$ and returns either the maximum or minimum value as appropriate as a $\langle\text{dimension denotation}\rangle$. For example

$\backslash\text{DimMathMax}$ $\{3.5\text{pt}\}$ $\{-2.7\text{pt}\}$	3.5pt
$\backslash\text{DimMathMin}$ $\{3.5\text{pt}\}$ $\{-2.7\text{pt}\}$	-2.7pt

9.3 Creating and Using Dimensions

 $\backslash\text{DimNew}$ $\langle\text{dimension}\rangle$

Creates a new $\langle\text{dimension}\rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle\text{dimension}\rangle$ is initially equal to 0 pt.

\DimConst $\langle dimension \rangle$ { $\langle dimension expression \rangle$ }

Creates a new constant $\langle dimension \rangle$ or raises an error if the name is already taken. The value of the $\langle dimension \rangle$ is set globally to the $\langle dimension expression \rangle$. For example

```
\DimConst \cFooSomeDim {1cm}
\DimUse \cFooSomeDim
```

28.45274pt

\DimUse $\langle dimension \rangle$

Recovers the content of a $\langle dimension \rangle$ and returns the value. An error is raised if the variable does not exist or if it is invalid.

9.4 Viewing Dimensions

\DimLog { $\langle dimension expression \rangle$ }

Writes the result of evaluating the $\langle dimension expression \rangle$ in the log file. For example

```
\DimLog {\lFooSomeDim+1cm}
```

\DimVarLog $\langle dimension \rangle$

Writes the value of the $\langle dimension \rangle$ in the log file. For example

```
\DimVarLog \lFooSomeDim
```

\DimShow { $\langle dimension expression \rangle$ }

Displays the result of evaluating the $\langle dimension expression \rangle$ on the terminal. For example

```
\DimShow {\lFooSomeDim+1cm}
```

\DimVarShow $\langle dimension \rangle$

Displays the value of the $\langle dimension \rangle$ on the terminal. For example

```
\DimVarShow \lFooSomeDim
```

9.5 Setting Dimension Variables

\DimSet $\langle dimension \rangle$ { $\langle dimension expression \rangle$ }

Sets $\langle dimension \rangle$ to the value of $\langle dimension expression \rangle$, which must evaluate to a length with units.

\DimSetEq $\langle dimension_1 \rangle$ $\langle dimension_2 \rangle$

Sets the content of $\langle dimension_1 \rangle$ equal to that of $\langle dimension_2 \rangle$. For example

```
\DimSet \lTmpaDim {10pt}
\DimSetEq \lTmpbDim \lTmpaDim
\DimUse \lTmpbDim
```

10.0pt

\DimZero *<dimension>*

Sets *<dimension>* to 0 pt. For example

```
\DimSet \lTmpaDim {1em}
\DimZero \lTmpaDim
\DimUse \lTmpaDim
```

0.0pt

\DimZeroNew *<dimension>*

Ensures that the *<dimension>* exists globally by applying **\DimNew** if necessary, then applies **\DimZero** to set the *<dimension>* to zero. For example

```
\DimZeroNew \lFooSomeDim
\DimUse \lFooSomeDim
```

0.0pt

\DimAdd *<dimension>* {*<dimension expression>*}

Adds the result of the *<dimension expression>* to the current content of the *<dimension>*. For example

```
\DimSet \lTmpaDim {5.3pt}
\DimAdd \lTmpaDim {2.11pt}
\DimUse \lTmpaDim
```

7.41pt

\DimSub *<dimension>* {*<dimension expression>*}

Subtracts the result of the *<dimension expression>* from the current content of the *<dimension>*. For example

```
\DimSet \lTmpaDim {5.3pt}
\DimSub \lTmpaDim {2.11pt}
\DimUse \lTmpaDim
```

3.19pt

9.6 Dimension Step Functions

\DimStepInline {*<initial value>*} {*<step>*} {*<final value>*} {*<code>*}

This function first evaluates the *<initial value>*, *<step>* and *<final value>*, all of which should be dimension expressions. Then for each *<value>* from the *<initial value>* to the *<final value>* in turn (using *<step>* between each *<value>*), the *<code>* is inserted into the input stream with #1 replaced by the current *<value>*. Thus the *<code>* should define a function of one argument (#1). For example

```

\IgnoreSpacesOn
\TlClear \lTmptl
\DimStepInline {1pt} {0.1pt} {1.5pt} {
  \TlPutRight \lTmptl {[#1]}
}
\Result {\Value\lTmptl}
\IgnoreSpacesOff

```

produces [1.0pt][1.1pt][1.20001pt][1.30002pt][1.40002pt].

\DimStepVariable {<initial value>} {<step>} {<final value>} <tl var> {<code>}

This function first evaluates the <initial value>, <step> and <final value>, all of which should be dimension expressions. Then for each <value> from the <initial value> to the <final value> in turn (using <step> between each <value>), the <code> is inserted into the input stream, with the <tl var> defined as the current <value>. Thus the <code> should make use of the <tl var>.

9.7 Dimension Conditionals

\DimIfExist <dimension>
\DimIfExistTF <dimension> {<true code>} {<false code>}

Tests whether the <dimension> is currently defined. This does not check that the <dimension> really is a dimension variable. For example

```

\DimIfExistTF \lTmptDim {\Result{Yes}} {\Result{No}}
\DimIfExistTF \lFooUndefinedDim {\Result{Yes}} {\Result{No}}

```

Yes No

\DimCompare {<dimexpr₁>} <relation> {<dimexpr₂>}
\DimCompareTF {<dimexpr₁>} <relation> {<dimexpr₂>} {<true code>} {<false code>}

This function first evaluates each of the <dimension expressions> as described for **\DimEval**. The two results are then compared using the <relation>:

Equal	=
Greater than	>
Less than	<

For example

```

\DimCompareTF {1pt} > {0.9999pt} {\Result{Greater}} {\Result{Less}}
\DimCompareTF {1pt} > {1.0001pt} {\Result{Greater}} {\Result{Less}}

```

Greater Less

9.8 Dimension Case Functions

```
\DimCase {⟨test dimension expression⟩}
{
  {⟨dimexpr case1⟩} {⟨code case1⟩}
  {⟨dimexpr case2⟩} {⟨code case2⟩}
  ...
  {⟨dimexpr casen⟩} {⟨code casen⟩}
}
```

This function evaluates the *⟨test dimension expression⟩* and compares this in turn to each of the *⟨dimension expression cases⟩*. If the two are equal then the associated *⟨code⟩* is left in the input stream and other cases are discarded.

```
\DimCaseT {⟨test dimension expression⟩}
{
  {⟨dimexpr case1⟩} {⟨code case1⟩}
  {⟨dimexpr case2⟩} {⟨code case2⟩}
  ...
  {⟨dimexpr casen⟩} {⟨code casen⟩}
}
{⟨true code⟩}
```

This function evaluates the *⟨test dimension expression⟩* and compares this in turn to each of the *⟨dimension expression cases⟩*. If the two are equal then the associated *⟨code⟩* is left in the input stream and other cases are discarded. If any of the cases are matched, the *⟨true code⟩* is also inserted into the input stream (after the code for the appropriate case).

```
\DimCaseF {⟨test dimension expression⟩}
{
  {⟨dimexpr case1⟩} {⟨code case1⟩}
  {⟨dimexpr case2⟩} {⟨code case2⟩}
  ...
  {⟨dimexpr casen⟩} {⟨code casen⟩}
}
{⟨false code⟩}
```

This function evaluates the *⟨test dimension expression⟩* and compares this in turn to each of the *⟨dimension expression cases⟩*. If the two are equal then the associated *⟨code⟩* is left in the input stream and other cases are discarded. If none of the cases match then the *⟨false code⟩* is inserted. For example

```
\IgnoreSpacesOn
\DimSet \lTmpaDim {5pt}
\DimCaseF {2\lTmpaDim} {
  {5pt}      {\Result{Small}}
  {4pt+6pt}  {\Result{Medium}}
  {-10pt}    {\Result{Negative}}
}{
  \Result {No Match}
}
\IgnoreSpacesOff
```

Medium

```

\DimCaseTF {⟨test dimension expression⟩}
{
  {⟨dimexpr case1⟩} {⟨code case1⟩}
  {⟨dimexpr case2⟩} {⟨code case2⟩}
  ...
  {⟨dimexpr casen⟩} {⟨code casen⟩}
}
{⟨true code⟩}
{⟨false code⟩}

```

This function evaluates the $\langle test\ dimension\ expression \rangle$ and compares this in turn to each of the $\langle dimension\ expression\ cases \rangle$. If the two are equal then the associated $\langle code \rangle$ is left in the input stream and other cases are discarded. If any of the cases are matched, the $\langle true\ code \rangle$ is also inserted into the input stream (after the code for the appropriate case), while if none match then the $\langle false\ code \rangle$ is inserted.

Chapter 10

Comma Separated Lists (Clist)

10.1 Constant and Scratch Comma Lists

`\cEmptyClist`

Constant that is always empty.

`\lTmpaClist \lTmpbClist \lTmpcClist \lTmpiClist \lTnpjClist \lTmpkClist`

Scratch comma lists for local assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

`\gTmpaClist \gTmpbClist \gTmpcClist \gTmpiClist \gTnpjClist \gTmpkClist`

Scratch comma lists for global assignment. These are never used by the **functional** package, and so are safe for use with any function. However, they may be overwritten by other code and so should only be used for short-term storage.

10.2 Creating and Using Comma Lists

`\ClistNew` $\langle comma\ list \rangle$

Creates a new $\langle comma\ list \rangle$ or raises an error if the name is already taken. The declaration is global. The $\langle comma\ list \rangle$ initially contains no items.

`\ClistConst` $\langle clist\ var \rangle$ $\{ \langle comma\ list \rangle \}$

Creates a new constant $\langle clist\ var \rangle$ or raises an error if the name is already taken. The value of the $\langle clist\ var \rangle$ is set globally to the $\langle comma\ list \rangle$.

`\ClistVarJoin` $\langle clist\ var \rangle$ $\{ \langle separator \rangle \}$

Returns the contents of the $\langle clist\ var \rangle$, with the $\langle separator \rangle$ between the items. For example,

```
\ClistSet \lTmpaClist { a , b , , c , {de} , f }
\ClistVarJoin \lTmpaClist { and }
```

a and b and c and de and f

\ClistVarJoinExtended $\langle clist\ var \rangle$ $\{\langle separator\ between\ two \rangle\}$ $\{\langle separator\ between\ more\ than\ two \rangle\}$ $\{\langle separator\ between\ final\ two \rangle\}$

Returns the contents of the $\langle clist\ var \rangle$, with the appropriate $\langle separator \rangle$ between the items. Namely, if the comma list has more than two items, the $\langle separator\ between\ more\ than\ two \rangle$ is placed between each pair of items except the last, for which the $\langle separator\ between\ final\ two \rangle$ is used. If the comma list has exactly two items, then they are joined with the $\langle separator\ between\ two \rangle$ and returns. For example,

```
\ClistSet \lTmPaClist { a , b , , c , {de} , f }
\ClistVarJoinExtended \lTmPaClist { and } {, } {, and }
```

a, b, c, de, and f

The first separator argument is not used in this case because the comma list has more than 2 items.

\ClistJoin $\langle comma\ list \rangle$ $\{\langle separator \rangle\}$
\ClistJoinExtended $\langle comma\ list \rangle$ $\{\langle separator\ between\ two \rangle\}$ $\{\langle separator\ between\ more\ than\ two \rangle\}$ $\{\langle separator\ between\ final\ two \rangle\}$

Returns the contents of the $\langle comma\ list \rangle$, with the appropriate $\langle separator \rangle$ between the items. As for **\ClistSet**, blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. The $\langle separators \rangle$ are then inserted in the same way as for **\ClistVarJoin** and **\ClistVarJoinExtended**, respectively.

10.3 Viewing Comma Lists

\ClistLog $\{\langle tokens \rangle\}$

Writes the entries in the comma list in the log file. See also **\ClistShow** which displays the result in the terminal.

\ClistVarLog $\langle comma\ list \rangle$

Writes the entries in the $\langle comma\ list \rangle$ in the log file. See also **\ClistVarShow** which displays the result in the terminal.

\ClistShow $\{\langle tokens \rangle\}$

Displays the entries in the comma list in the terminal.

\ClistVarShow $\langle comma\ list \rangle$

Displays the entries in the $\langle comma\ list \rangle$ in the terminal.

10.4 Setting Comma Lists

\ClistSet $\langle comma\ list \rangle$ $\{\langle item_1 \rangle, ..., \langle item_n \rangle\}$

Sets $\langle comma\ list \rangle$ to contain the $\langle items \rangle$, removing any previous content from the variable. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To store some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the $\langle tokens \rangle$ contain commas or spaces, add a set of braces: **\ClistSet** $\langle comma\ list \rangle$ $\{\{\langle tokens \rangle\}\}$.

\ClistSetEq $\langle comma list_1 \rangle$ $\langle comma list_2 \rangle$

Sets the content of $\langle comma list_1 \rangle$ equal to that of $\langle comma list_2 \rangle$. To set a token list variable equal to a comma list variable, use **\TlSetEq**. Conversely, setting a comma list variable to a token list is unadvisable unless one checks space-trimming and related issues.

\ClistClear $\langle comma list \rangle$

Clears all items from the $\langle comma list \rangle$.

\ClistClearNew $\langle comma list \rangle$

Ensures that the $\langle comma list \rangle$ exists globally by applying **\ClistNew** if necessary, then applies **\ClistClear** to leave the list empty.

\ClistConcat $\langle comma list_1 \rangle$ $\langle comma list_2 \rangle$ $\langle comma list_3 \rangle$

Concatenates the content of $\langle comma list_2 \rangle$ and $\langle comma list_3 \rangle$ together and saves the result in $\langle comma list_1 \rangle$. The items in $\langle comma list_2 \rangle$ are placed at the left side of the new comma list.

\ClistPutLeft $\langle comma list \rangle$ $\{ \langle item_1 \rangle, \dots, \langle item_n \rangle \}$

Appends the $\langle items \rangle$ to the left of the $\langle comma list \rangle$. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To append some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the $\langle tokens \rangle$ contain commas or spaces, add a set of braces: **\ClistPutLeft** $\langle comma list \rangle$ $\{ \{ \langle tokens \rangle \} \}$.

\ClistPutRight $\langle comma list \rangle$ $\{ \langle item_1 \rangle, \dots, \langle item_n \rangle \}$

Appends the $\langle items \rangle$ to the right of the $\langle comma list \rangle$. Blank items are omitted, spaces are removed from both sides of each item, then a set of braces is removed if the resulting space-trimmed item is braced. To append some $\langle tokens \rangle$ as a single $\langle item \rangle$ even if the $\langle tokens \rangle$ contain commas or spaces, add a set of braces: **\ClistPutRight** $\langle comma list \rangle$ $\{ \{ \langle tokens \rangle \} \}$.

10.5 Modifying Comma Lists

While comma lists are normally used as ordered lists, it may be necessary to modify the content. The functions here may be used to update comma lists, while retaining the order of the unaffected entries.

\ClistRemoveDuplicates $\langle comma list \rangle$

Removes duplicate items from the $\langle comma list \rangle$, leaving the left most copy of each item in the $\langle comma list \rangle$. The $\langle item \rangle$ comparison takes place on a token basis, as for **\TlIfEqTF**. This function iterates through every item in the $\langle comma list \rangle$ and does a comparison with the $\langle items \rangle$ already checked. It is therefore relatively slow with large comma lists. Furthermore, it may fail if any of the items in the $\langle comma list \rangle$ contains $\{$, $\}$, or $\#$ (assuming the usual T_EX category codes apply).

\ClistRemoveAll $\langle comma list \rangle$ $\{ \langle item \rangle \}$

Removes every occurrence of $\langle item \rangle$ from the $\langle comma list \rangle$. The $\langle item \rangle$ comparison takes place on a token basis, as for **\TlIfEqTF**. The function may fail if the $\langle item \rangle$ contains $\{$, $\}$, or $\#$ (assuming the usual T_EX category codes apply).

\ClistReverse $\langle comma list \rangle$

Reverses the order of items stored in the $\langle comma list \rangle$.

10.6 Working with the Contents of Comma Lists

\ClistCount $\{\langle comma list \rangle\}$

\ClistVarCount $\langle comma list \rangle$

Returns the number of items in the $\langle comma list \rangle$ as an $\langle integer denotation \rangle$. The total number of items in a $\langle comma list \rangle$ includes those which are duplicates, *i.e.* every item in a $\langle comma list \rangle$ is counted.

\ClistItem $\{\langle comma list \rangle\} \{\langle integer expression \rangle\}$

Indexing items in the $\langle comma list \rangle$ from 1 at the top (left), this function evaluates the $\langle integer expression \rangle$ and returns the appropriate item from the comma list. If the $\langle integer expression \rangle$ is negative, indexing occurs from the bottom (right) of the comma list. When the $\langle integer expression \rangle$ is larger than the number of items in the $\langle comma list \rangle$ (as calculated by **\ClistCount**) then the function returns nothing.

\ClistVarItem $\langle comma list \rangle \{\langle integer expression \rangle\}$

Indexing items in the $\langle comma list \rangle$ from 1 at the top (left), this function evaluates the $\langle integer expression \rangle$ and returns the appropriate item from the comma list. If the $\langle integer expression \rangle$ is negative, indexing occurs from the bottom (right) of the comma list. When the $\langle integer expression \rangle$ is larger than the number of items in the $\langle comma list \rangle$ (as calculated by **\ClistVarCount**) then the function returns nothing.

\ClistRandItem $\{\langle comma list \rangle\}$

\ClistVarRandItem $\langle clist var \rangle$

Selects a pseudo-random item of the $\langle comma list \rangle$. If the $\langle comma list \rangle$ has no item, the result is empty.

10.7 Comma Lists as Stacks

Comma lists can be used as stacks, where data is pushed to and popped from the top of the comma list. (The left of a comma list is the top, for performance reasons.) The stack functions for comma lists are not intended to be mixed with the general ordered data functions detailed in the previous section: a comma list should either be used as an ordered data type or as a stack, but not in both ways.

\ClistGet $\langle comma list \rangle \langle token list variable \rangle$

\ClistGetT $\langle comma list \rangle \langle token list variable \rangle \langle true code \rangle$

\ClistGetF $\langle comma list \rangle \langle token list variable \rangle \langle false code \rangle$

\ClistGetTF $\langle comma list \rangle \langle token list variable \rangle \langle true code \rangle \langle false code \rangle$

Stores the left-most item from a $\langle comma list \rangle$ in the $\langle token list variable \rangle$ without removing it from the $\langle comma list \rangle$. The $\langle token list variable \rangle$ is assigned locally.

\ClistPop $\langle comma list \rangle \langle token list variable \rangle$

\ClistPopT $\langle comma list \rangle \langle token list variable \rangle \{\langle true code \rangle\}$

\ClistPopF $\langle comma list \rangle \langle token list variable \rangle \{\langle false code \rangle\}$

\ClistPopTF $\langle comma list \rangle \langle token list variable \rangle \{\langle true code \rangle\} \{\langle false code \rangle\}$

Pops the left-most item from a $\langle comma list \rangle$ into the $\langle token list variable \rangle$, *i.e.* removes the item from the comma list and stores it in the $\langle token list variable \rangle$. The assignment of the $\langle token list variable \rangle$ is local.

If the $\langle comma list \rangle$ is empty, the value of the $\langle token list variable \rangle$ is not defined in this case and should not be relied upon.

\ClistPush $\langle comma list \rangle$ $\{\langle items \rangle\}$

Adds the $\{\langle items \rangle\}$ to the top of the $\langle comma list \rangle$. Spaces are removed from both sides of each item as for any n-type comma list.

10.8 Mapping over Comma Lists

When the comma list is given explicitly, spaces are trimmed around each item. If the result of trimming spaces is empty, the item is ignored. Otherwise, if the item is surrounded by braces, one set is removed, and the result is passed to the mapped function. Thus, if the comma list that is being mapped is $\{a, \{b\}, \{c\}, \}$ then the arguments passed to the mapped function are ‘a’, ‘{b}’, an empty argument, and ‘c’.

When the comma list is given as a variable, spaces have already been trimmed on input, and items are simply stripped of one set of braces if any. This case is more efficient than using explicit comma lists.

\ClistMapInline $\{\langle comma list \rangle\}$ $\{\langle inline function \rangle\}$
\ClistVarMapInline $\langle comma list \rangle$ $\{\langle inline function \rangle\}$

Applies $\langle inline function \rangle$ to every $\langle item \rangle$ stored within the $\langle comma list \rangle$. The $\langle inline function \rangle$ should consist of code which receives the $\langle item \rangle$ as #1. The $\langle items \rangle$ are returned from left to right. For example

```
\IgnoreSpacesOn
\TlClear \lTmpl
\ClistMapInline {one,two,three} {
  \TlPutRight \lTmpl {(\#1)}
}
\Result {\TlUse\lTmpl}
\IgnoreSpacesOff
```

produces (one)(two)(three).

\ClistMapVariable $\{\langle comma list \rangle\}$ $\langle variable \rangle$ $\{\langle code \rangle\}$
\ClistVarMapVariable $\langle comma list \rangle$ $\langle variable \rangle$ $\{\langle code \rangle\}$

Stores each $\langle item \rangle$ of the $\langle comma list \rangle$ in turn in the (token list) $\langle variable \rangle$ and applies the $\langle code \rangle$. The $\langle code \rangle$ will usually make use of the $\langle variable \rangle$, but this is not enforced. The assignments to the $\langle variable \rangle$ are local. Its value after the loop is the last $\langle item \rangle$ in the $\langle comma list \rangle$, or its original value if there were no $\langle item \rangle$. The $\langle items \rangle$ are returned from left to right.

10.9 Comma List Conditionals

\ClistIfExist $\langle comma list \rangle$
\ClistIfExistTF $\langle comma list \rangle$ $\{\langle true code \rangle\}$ $\{\langle false code \rangle\}$

Tests whether the $\langle comma list \rangle$ is currently defined. This does not check that the $\langle comma list \rangle$ really is a comma list.

\ClistIfEmpty $\{\langle comma list \rangle\}$
\ClistIfEmptyTF $\{\langle comma list \rangle\}$ $\{\langle true code \rangle\}$ $\{\langle false code \rangle\}$

Tests if the $\langle comma list \rangle$ is empty (containing no items). The rules for space trimming are as for other

n-type comma-list functions, hence the comma list { , , , } (without outer braces) is empty, while { , { } , } (without outer braces) contains one element, which happens to be empty: the comma-list is not empty.

```
\ClistVarIfEmpty <comma list>
\ClistVarIfEmptyTF <comma list> {<true code>} {<false code>}
```

Tests if the *<comma list>* is empty (containing no items).

```
\ClistIfIn {<comma list>} {<item>}
\ClistIfInTF {<comma list>} {<item>} {<true code>} {<false code>}
```

Tests if the *<item>* is present in the *<comma list>*. In the case of an **n**-type *<comma list>*, the usual rules of space trimming and brace stripping apply. For example

```
\ClistIfInTF { a , {b} , {b} , c } {b} {Yes} {No}
```

Yes

```
\ClistVarIfIn <comma list> {<item>}
\ClistVarIfInTF <comma list> {<item>} {<true code>} {<false code>}
```

Tests if the *<item>* is present in the *<comma list>*. In the case of an **n**-type *<comma list>*, the usual rules of space trimming and brace stripping apply.

Chapter 11

The Source Code

```
%% -----
%% Functional: LaTeX2 functional interfaces for LaTeX3 programming layer
%% Copyright : 2022 (c) Jianrui Lyu <tolvjvr@163.com>
%% Repository: https://github.com/lvjvr/functional
%% Repository: https://bitbucket.org/lvjvr/functional
%% License   : The LaTeX Project Public License 1.3c
%% -----
```

11.1 Interfaces for Functional Programming (Prg)

```
\NeedsTeXFormat{LaTeX2e}[2018-04-01]

\RequirePackage{expl3}
\ProvidesExplPackage{functional}{2022-03-19}{2022B}
  {^^JLaTeX2 functional interfaces for LaTeX3 programming layer}

\cs_generate_variant:Nn \iow_log:n { V }
\cs_generate_variant:Nn \tl_log:n { e }
\cs_generate_variant:Nn \tl_set:Nn { Ne }

\tl_new:N \gResultTl
\int_new:N \l__fun_arg_count_int
\tl_new:N \l__fun_parameters_defined_tl
\tl_const:Nn \c__fun_parameter_defined_i__tl { } % no argument
\tl_const:Nn \c__fun_parameter_defined_i_i_tl { #1 }
\tl_const:Nn \c__fun_parameter_defined_i_ii_tl { #1 #2 }
\tl_const:Nn \c__fun_parameter_defined_i_iii_tl { #1 #2 #3 }
\tl_const:Nn \c__fun_parameter_defined_i_iv_tl { #1 #2 #3 #4 }
\tl_const:Nn \c__fun_parameter_defined_i_v_tl { #1 #2 #3 #4 #5 }
\tl_const:Nn \c__fun_parameter_defined_i_vi_tl { #1 #2 #3 #4 #5 #6 }
\tl_const:Nn \c__fun_parameter_defined_i_vii_tl { #1 #2 #3 #4 #5 #6 #7 }
\tl_const:Nn \c__fun_parameter_defined_i_viii_tl { #1 #2 #3 #4 #5 #6 #7 #8 }
\tl_const:Nn \c__fun_parameter_defined_i_ix_tl { #1 #2 #3 #4 #5 #6 #7 #8 #9 }
\tl_new:N \l__fun_parameters_called_tl
\tl_const:Nn \c__fun_parameter_called_i_i_tl { {#1} }
\tl_const:Nn \c__fun_parameter_called_i_ii_tl { {#1}{#2} }
\tl_const:Nn \c__fun_parameter_called_i_iii_tl { {#1}{#2}{#3} }
\tl_const:Nn \c__fun_parameter_called_i_iv_tl { {#1}{#2}{#3}{#4} }
\tl_const:Nn \c__fun_parameter_called_i_v_tl { {#1}{#2}{#3}{#4}{#5} }
\tl_const:Nn \c__fun_parameter_called_i_vi_tl { {#1}{#2}{#3}{#4}{#5}{#6} }
\tl_const:Nn \c__fun_parameter_called_i_vii_tl { {#1}{#2}{#3}{#4}{#5}{#6}{#7} }
```

```

\tl_new:N \l__fun_parameters_true_tl
\tl_new:N \l__fun_parameters_false_tl
\tl_const:Nn \c__fun_parameter_called_i_tl { {#1} }
\tl_const:Nn \c__fun_parameter_called_ii_tl { {#2} }
\tl_const:Nn \c__fun_parameter_called_iii_tl { {#3} }
\tl_const:Nn \c__fun_parameter_called_iv_tl { {#4} }
\tl_const:Nn \c__fun_parameter_called_v_tl { {#5} }
\tl_const:Nn \c__fun_parameter_called_vi_tl { {#6} }
\tl_const:Nn \c__fun_parameter_called_vii_tl { {#7} }
\tl_const:Nn \c__fun_parameter_called_viii_tl { {#8} }
\tl_const:Nn \c__fun_parameter_called_ix_tl { {#9} }

%% #1: function name; #2: argument specification; #3 function body
\cs_new_protected:Npn \__fun_new_function:Nnn #1 #2 #3
{
  \int_set:Nn \l__fun_arg_count_int { \tl_count:n {#2} } % spaces are ignored
  \tl_set_eq:Nc \l__fun_parameters_defined_tl
    { c__fun_parameter_defined_i_ \int_to_roman:n { \l__fun_arg_count_int } _tl }
  \exp_last_unbraced:NcV \cs_new_protected:Npn
    { __fun_defined_ \cs_to_str:N #1 : w }
    \l__fun_parameters_defined_tl
    {
      \__fun_group_begin:
      \tl_gclear:N \gResultTl
      #3
      \__fun_tracing_log:e { [0] ~ \gResultTl }
      \__fun_group_end:
    }
  \use:c { __fun_new_with_arg_ \int_to_roman:n { \l__fun_arg_count_int } :NnV }
    #1 {#2} \l__fun_parameters_defined_tl
}
\cs_generate_variant:Nn \__fun_new_function:Nnn { cne }

\cs_set_eq:NN \PrgNewFunction \__fun_new_function:Nnn

\tl_new:N \g__fun_last_result_tl

%% #1: function name; #2: argument specification; #3 function body
\cs_new_protected:Npn \__fun_new_conditional:Nnn #1 #2 #3
{
  \__fun_new_function:Nnn #1 { #2 } { #3 }
  \tl_set_eq:Nc \l__fun_parameters_called_tl
    { c__fun_parameter_called_i_ \int_to_roman:n { \l__fun_arg_count_int } _tl }
  \tl_set_eq:Nc \l__fun_parameters_true_tl
    { c__fun_parameter_called_ \int_to_roman:n { \l__fun_arg_count_int + 1 } _tl }
  \tl_set_eq:Nc \l__fun_parameters_false_tl
    { c__fun_parameter_called_ \int_to_roman:n { \l__fun_arg_count_int + 2 } _tl }
  \__fun_new_function:cne { \cs_to_str:N #1 TF } { #2 n n }
  {
    #1 \exp_not:V \l__fun_parameters_called_tl
    \exp_not:n
    {
      \tl_set_eq:NN \g__fun_last_result_tl \gResultTl
      \tl_gclear:N \gResultTl
      \exp_last_unbraced:NV \bool_if:NTF \g__fun_last_result_tl
    }
    \exp_not:V \l__fun_parameters_true_tl
  }
}

```



```

        \exp_not:V \l__fun_parameters_false_tl
    }
}

\cs_set_eq:NN \PrgNewConditional \__fun_new_conditional:Nnn

\int_new:N \g__fun_nesting_level_int

%% #1: function name; #2: argument specifications; #3 parameters tl defined
%% Some times we need to create a function without arguments
\cs_new_protected:Npn \__fun_new_with_arg:Nnn #1 #2 #3
{
    \cs_new_protected:Npn #1 #3
    {
        \int_gincr:N \g__fun_nesting_level_int
        \__fun_evaluate:Nn #1 {#2}
        \int_gdecr:N \g__fun_nesting_level_int
        \__fun_return_result:
    }
}

\cs_generate_variant:Nn \__fun_new_with_arg:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_i:Nnn #1 #2 #3
{
    \cs_new_protected:Npn #1 #3
    {
        \int_gincr:N \g__fun_nesting_level_int
        \__fun_one_argument_gset:nn { 1 } { ##1 }
        \__fun_evaluate:Nn #1 {#2}
        \int_gdecr:N \g__fun_nesting_level_int
        \__fun_return_result:
    }
}

\cs_generate_variant:Nn \__fun_new_with_arg_i:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_ii:Nnn #1 #2 #3
{
    \cs_new_protected:Npn #1 #3
    {
        \int_gincr:N \g__fun_nesting_level_int
        \__fun_one_argument_gset:nn { 1 } { ##1 }
        \__fun_one_argument_gset:nn { 2 } { ##2 }
        \__fun_evaluate:Nn #1 {#2}
        \int_gdecr:N \g__fun_nesting_level_int
        \__fun_return_result:
    }
}

\cs_generate_variant:Nn \__fun_new_with_arg_ii:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_iii:Nnn #1 #2 #3
{
    \cs_new_protected:Npn #1 #3
    {
        \int_gincr:N \g__fun_nesting_level_int

```

```

    \__fun_one_argument_gset:nn { 1 } { ##1 }
    \__fun_one_argument_gset:nn { 2 } { ##2 }
    \__fun_one_argument_gset:nn { 3 } { ##3 }
    \__fun_evaluate:Nn #1 {#2}
    \int_gdecr:N \g__fun_nesting_level_int
    \__fun_return_result:
  }
}
\cs_generate_variant:Nn \__fun_new_with_arg_iii:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_iv:Nnn #1 #2 #3
{
  \cs_new_protected:Npn #1 #3
  {
    \int_gincr:N \g__fun_nesting_level_int
    \__fun_one_argument_gset:nn { 1 } { ##1 }
    \__fun_one_argument_gset:nn { 2 } { ##2 }
    \__fun_one_argument_gset:nn { 3 } { ##3 }
    \__fun_one_argument_gset:nn { 4 } { ##4 }
    \__fun_evaluate:Nn #1 {#2}
    \int_gdecr:N \g__fun_nesting_level_int
    \__fun_return_result:
  }
}
\cs_generate_variant:Nn \__fun_new_with_arg_iv:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_v:Nnn #1 #2 #3
{
  \cs_new_protected:Npn #1 #3
  {
    \int_gincr:N \g__fun_nesting_level_int
    \__fun_one_argument_gset:nn { 1 } { ##1 }
    \__fun_one_argument_gset:nn { 2 } { ##2 }
    \__fun_one_argument_gset:nn { 3 } { ##3 }
    \__fun_one_argument_gset:nn { 4 } { ##4 }
    \__fun_one_argument_gset:nn { 5 } { ##5 }
    \__fun_evaluate:Nn #1 {#2}
    \int_gdecr:N \g__fun_nesting_level_int
    \__fun_return_result:
  }
}
\cs_generate_variant:Nn \__fun_new_with_arg_v:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_vi:Nnn #1 #2 #3
{
  \cs_new_protected:Npn #1 #3
  {
    \int_gincr:N \g__fun_nesting_level_int
    \__fun_one_argument_gset:nn { 1 } { ##1 }
    \__fun_one_argument_gset:nn { 2 } { ##2 }
    \__fun_one_argument_gset:nn { 3 } { ##3 }
    \__fun_one_argument_gset:nn { 4 } { ##4 }
    \__fun_one_argument_gset:nn { 5 } { ##5 }
    \__fun_one_argument_gset:nn { 6 } { ##6 }
    \__fun_evaluate:Nn #1 {#2}
  }
}

```

```

        \int_gdecr:N \g__fun_nesting_level_int
        \__fun_return_result:
    }
}
\cs_generate_variant:Nn \__fun_new_with_arg_vi:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_vii:Nnn #1 #2 #3
{
    \cs_new_protected:Npn #1 #3
    {
        \int_gincr:N \g__fun_nesting_level_int
        \__fun_one_argument_gset:nn { 1 } { ##1 }
        \__fun_one_argument_gset:nn { 2 } { ##2 }
        \__fun_one_argument_gset:nn { 3 } { ##3 }
        \__fun_one_argument_gset:nn { 4 } { ##4 }
        \__fun_one_argument_gset:nn { 5 } { ##5 }
        \__fun_one_argument_gset:nn { 6 } { ##6 }
        \__fun_one_argument_gset:nn { 7 } { ##7 }
        \__fun_evaluate:Nn #1 {#2}
        \int_gdecr:N \g__fun_nesting_level_int
        \__fun_return_result:
    }
}
\cs_generate_variant:Nn \__fun_new_with_arg_vii:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_viii:Nnn #1 #2 #3
{
    \cs_new_protected:Npn #1 #3
    {
        \int_gincr:N \g__fun_nesting_level_int
        \__fun_one_argument_gset:nn { 1 } { ##1 }
        \__fun_one_argument_gset:nn { 2 } { ##2 }
        \__fun_one_argument_gset:nn { 3 } { ##3 }
        \__fun_one_argument_gset:nn { 4 } { ##4 }
        \__fun_one_argument_gset:nn { 5 } { ##5 }
        \__fun_one_argument_gset:nn { 6 } { ##6 }
        \__fun_one_argument_gset:nn { 7 } { ##7 }
        \__fun_one_argument_gset:nn { 8 } { ##8 }
        \__fun_evaluate:Nn #1 {#2}
        \int_gdecr:N \g__fun_nesting_level_int
        \__fun_return_result:
    }
}
\cs_generate_variant:Nn \__fun_new_with_arg_viii:Nnn { NnV }

%% #1: function name; #2: argument specifications; #3 parameters tl defined
\cs_new_protected:Npn \__fun_new_with_arg_ix:Nnn #1 #2 #3
{
    \cs_new_protected:Npn #1 #3
    {
        \int_gincr:N \g__fun_nesting_level_int
        \__fun_one_argument_gset:nn { 1 } { ##1 }
        \__fun_one_argument_gset:nn { 2 } { ##2 }
        \__fun_one_argument_gset:nn { 3 } { ##3 }
        \__fun_one_argument_gset:nn { 4 } { ##4 }
        \__fun_one_argument_gset:nn { 5 } { ##5 }
    }
}

```

```

    \__fun_one_argument_gset:nn { 6 } { ##6 }
    \__fun_one_argument_gset:nn { 7 } { ##7 }
    \__fun_one_argument_gset:nn { 8 } { ##8 }
    \__fun_one_argument_gset:nn { 9 } { ##9 }
    \__fun_evaluate:Nn #1 {#2}
    \int_gdecr:N \g__fun_nesting_level_int
    \__fun_return_result:
  }
}
\cs_generate_variant:Nn \__fun_new_with_arg_ix:Nnn { NnV }

\tl_new:N \l__fun_argtype_tl
\tl_const:Nn \c__fun_argtype_m_tl { m }
\tl_const:Nn \c__fun_argtype_M_tl { M }
\tl_const:Nn \c__fun_argtype_n_tl { n }
\tl_const:Nn \c__fun_argtype_N_tl { N }
\tl_new:N \l__fun_argument_tl

%% #1: function name; #2: argument specifications
\cs_new_protected:Npn \__fun_evaluate:Nn #1 #2
{
  \__fun_argtype_index_gzero:
  \__fun_arguments_gclear:
  \tl_map_variable:nNn { #2 } \l__fun_argtype_tl % spaces are ignored
  {
    \__fun_argtype_index_gincr:
    \__fun_one_argument_get:eN { \__fun_argtype_index_use: } \l__fun_argument_tl
    \tl_case:Nn \l__fun_argtype_tl
    {
      \c__fun_argtype_m_tl
      {
        \__fun_evaluate_and_put_argument:N \l__fun_argument_tl
      }
      \c__fun_argtype_M_tl
      {
        \__fun_evaluate_and_put_argument:N \l__fun_argument_tl
      }
      \c__fun_argtype_n_tl
      {
        \__fun_arguments_gput:e { { \exp_not:V \l__fun_argument_tl } }
      }
      \c__fun_argtype_N_tl
      {
        \__fun_arguments_gput:e { \exp_not:V \l__fun_argument_tl }
      }
    }
  }
}

\__fun_arguments_log:N #1
\__fun_arguments_called:c { __fun_defined_ \cs_to_str:N #1 : w }
}

\cs_new_protected:Npn \__fun_evaluate_and_put_argument:N #1
{
  \cs_if_exist:cTF
  {
    __fun_defined_ \exp_last_unbraced:Ne \cs_to_str:N { \tl_head:N #1 } : w
  }
  {

```

```

        #1
        \__fun_arguments_gput:e { { \exp_not:V \gResultTl } }
    }
    {
        \__fun_arguments_gput:e { { \exp_not:V #1 } }
    }
}

%% #1: argument number; #2: token lists
\cs_new_protected:Npn \__fun_one_argument_gset:nn #1 #2
{
    \tl_gset:cn
    { g__fun_one_argument_ \int_use:N \g__fun_nesting_level_int _#1_tl } { #2 }
    %\__fun_one_argument_log:nn { #1 } { set }
}

%% #1: argument number; #2: variable of token lists
\cs_new_protected:Npn \__fun_one_argument_get:nN #1 #2
{
    \tl_set_eq:Nc
    #2 { g__fun_one_argument_ \int_use:N \g__fun_nesting_level_int _#1_tl }
    %\__fun_one_argument_log:nn { #1 } { get }
}
\cs_generate_variant:Nn \__fun_one_argument_get:nN { eN }

%% #1: argument number; #2: get or set
\cs_new_protected:Npn \__fun_one_argument_log:nn #1 #2
{
    \tl_log:e
    {
        #2 ~ level _ \int_use:N \g__fun_nesting_level_int _ arg _ #1 ~ = ~
        \exp_not:v
        { g__fun_one_argument_ \int_use:N \g__fun_nesting_level_int _#1_tl }
    }
}

\int_new:c { g__fun_argtype_index_ 1 _int }
\int_new:c { g__fun_argtype_index_ 2 _int }
\int_new:c { g__fun_argtype_index_ 3 _int }
\int_new:c { g__fun_argtype_index_ 4 _int }
\int_new:c { g__fun_argtype_index_ 5 _int }

\cs_new_protected:Npn \__fun_argtype_index_gzero:
{
    \int_gzero_new:c
    { g__fun_argtype_index_ \int_use:N \g__fun_nesting_level_int _int }
}

\cs_new_protected:Npn \__fun_argtype_index_gincr:
{
    \int_gincr:c
    { g__fun_argtype_index_ \int_use:N \g__fun_nesting_level_int _int }
}

\cs_new:Npn \__fun_argtype_index_use:
{

```

```

\int_use:c { g__fun_argtype_index_ \int_use:N \g__fun_nesting_level_int _int }
}

\cs_new_protected:Npn \__fun_arguments_called:N #1
{
  \exp_last_unbraced:Nv
    #1 { g__fun_arguments_ \int_use:N \g__fun_nesting_level_int _tl }
}
\cs_generate_variant:Nn \__fun_arguments_called:N { c }

\cs_new_protected:Npn \__fun_arguments_gclear:
{
  \tl_gclear:c { g__fun_arguments_ \int_use:N \g__fun_nesting_level_int _tl }
}

\cs_new_protected:Npn \__fun_arguments_log:N #1
{
  \__fun_tracing_log:e
  {
    [I] ~ \token_to_str:N #1
    \exp_not:v { g__fun_arguments_ \int_use:N \g__fun_nesting_level_int _tl }
  }
}

\cs_new_protected:Npn \__fun_arguments_gput:n #1
{
  \tl_gput_right:cn
    { g__fun_arguments_ \int_use:N \g__fun_nesting_level_int _tl } { #1 }
}
\cs_generate_variant:Nn \__fun_arguments_gput:n { e }

\cs_set_eq:NN \Break \prg_break:
\cs_set_eq:NN \PrgBreak \prg_break:

\cs_set_eq:NN \BreakDo \prg_break:n
\cs_set_eq:NN \PrgBreakDo \prg_break:n

\cs_new_protected:Npn \__fun_put_result:n #1
{
  \tl_gput_right:Nn \gResultTl { #1 }
}
\cs_generate_variant:Nn \__fun_put_result:n { e, V }

\PrgNewFunction \Result { m }
{
  \__fun_put_result:n { #1 }
}

\cs_new_protected:Npn \__fun_return_result:
{
  \int_compare:nNnT { \g__fun_nesting_level_int } = { 0 }
  { \tl_use:N \gResultTl }
}

\tl_new:N \l__fun_variable_type_tl

```

```

\prg_new_protected_conditional:Npnn \__fun_if_global_variable:N #1 { TF }
{
  \tl_set:Nc \l__fun_variable_type_tl
    { \exp_args:Nc \tl_head:n { \cs_to_str:N #1 } }
  \str_if_eq:VnTF \l__fun_variable_type_tl { g }
  { \prg_return_true: }
  {
    \str_if_eq:VnTF \l__fun_variable_type_tl { c }
    { \prg_return_true: }
    { \prg_return_false: }
  }
}

%% We must not put an assignment inside a group
\cs_new_protected:Npn \__fun_do_assignment:Nnn #1 #2 #3
{
  \__fun_group_end:
  \__fun_if_global_variable:NTF #1 { #2 } { #3 }
  \__fun_group_begin:
}

\bool_new:N \l__fun_scoping_bool

\cs_new_protected:Npn \__fun_scoping_true:
{
  \cs_set_eq:NN \__fun_group_begin: \group_begin:
  \cs_set_eq:NN \__fun_group_end: \group_end:
}

\cs_new_protected:Npn \__fun_scoping_false:
{
  \cs_set_eq:NN \__fun_group_begin: \scan_stop:
  \cs_set_eq:NN \__fun_group_end: \scan_stop:
}

\cs_new_protected:Npn \__fun_scoping_set:
{
  \bool_if:NTF \l__fun_scoping_bool
    { \__fun_scoping_true: } { \__fun_scoping_false: }
}

\bool_new:N \l__fun_tracing_bool
\tl_new:N \l__tracing_text_tl

\cs_new_protected:Npn \__fun_tracing_log_on:n #1
{
  \tl_set:Nc \l__tracing_text_tl
    {
      \prg_replicate:nn
        { \int_eval:n { (\g__fun_nesting_level_int - 1) * 4 } } { ~ }
    }
  \tl_put_right:Nn \l__tracing_text_tl { #1 }
  \iow_log:V \l__tracing_text_tl
}

\cs_generate_variant:Nn \__fun_tracing_log_on:n { e, V }

\cs_new_protected:Npn \__fun_tracing_log_off:n #1 { }

```

```

\cs_generate_variant:Nn \__fun_tracing_log_off:n { e, V }

\cs_new_protected:Npn \__fun_tracing_true:
{
  \cs_set_eq:NN \__fun_tracing_log:n \__fun_tracing_log_on:n
  \cs_set_eq:NN \__fun_tracing_log:e \__fun_tracing_log_on:e
  \cs_set_eq:NN \__fun_tracing_log:V \__fun_tracing_log_on:V
}

\cs_new_protected:Npn \__fun_tracing_false:
{
  \cs_set_eq:NN \__fun_tracing_log:n \__fun_tracing_log_off:n
  \cs_set_eq:NN \__fun_tracing_log:e \__fun_tracing_log_off:e
  \cs_set_eq:NN \__fun_tracing_log:V \__fun_tracing_log_off:V
}

\cs_new_protected:Npn \__fun_tracing_set:
{
  \bool_if:NTF \l__fun_tracing_bool
  { \__fun_tracing_true: } { \__fun_tracing_false: }
}

\keys_define:nn { functional }
{
  scoping .bool_set:N = \l__fun_scoping_bool,
  tracing .bool_set:N = \l__fun_tracing_bool,
}

\NewDocumentCommand \Functional { m }
{
  \keys_set:nn { functional } { #1 }
  \__fun_scoping_set:
  \__fun_tracing_set:
}

\Functional { scoping = false, tracing = false }

\cs_new_protected:Npn \__fun_ignore_spaces_on:
{
  \ExplSyntaxOn
  \char_set_catcode_math_subscript:N \_
  \char_set_catcode_other:N \:
}
\cs_set_eq:NN \IgnoreSpacesOn \__fun_ignore_spaces_on:
\cs_set_eq:NN \IgnoreSpacesOff \ExplSyntaxOff

```

11.2 Interfaces for Argument Using (Use)

```

\PrgNewFunction \Name { m }
{
  \exp_args:Nc \__fun_put_result:n { #1 }
}
\cs_set_eq:NN \UseName \Name

\PrgNewFunction \Value { M }

```



```

{
    __fun_put_result:V #1
}
\cs_set_eq:NN \UseValue \Value

\PrgNewFunction \Expand { m }
{
    __fun_put_result:e { #1 }
}
\cs_set_eq:NN \UseExpand \Expand

\cs_set_eq:NN \ExpNot \exp_not:n
\cs_set_eq:NN \ExpValue \exp_not:V

\PrgNewFunction \UseOne { n } { \Result { #1 } }

\PrgNewFunction \GobbleOne { n } { \Result { } }

\PrgNewFunction \UseGobble { n n } { \UseOne { #1 } }

\PrgNewFunction \GobbleUse { n n } { \UseOne { #2 } }

```

11.3 Interfaces for Control Structures (Bool)

```

\bool_const:Nn \cTrueBool { \c_true_bool }
\bool_const:Nn \cFalseBool { \c_false_bool }

\bool_new:N \lTmpaBool \bool_new:N \lTmpbBool \bool_new:N \lTmpcBool
\bool_new:N \lTmpiBool \bool_new:N \lTmpjBool \bool_new:N \lTmpkBool
\bool_new:N \l@Funx@Bool \bool_new:N \l@Funy@Bool \bool_new:N \l@Funz@Bool

\bool_new:N \gTmpaBool \bool_new:N \gTmpbBool \bool_new:N \gTmpcBool
\bool_new:N \gTmpiBool \bool_new:N \gTmpjBool \bool_new:N \gTmpkBool
\bool_new:N \g@Funx@Bool \bool_new:N \g@Funy@Bool \bool_new:N \g@Funz@Bool

\PrgNewFunction \BoolNew { M } { \bool_new:N #1 }

\PrgNewFunction \BoolConst { M m } { \bool_const:Nn #1 { #2 } }

\PrgNewFunction \BoolSet { M m } {
    __fun_do_assignment:Nnn #1
    { \bool_gset:Nn #1 { #2 } } { \bool_set:Nn #1 { #2 } }
}

\PrgNewFunction \BoolSetTrue { M }
{
    __fun_do_assignment:Nnn #1 { \bool_gset_true:N #1 } { \bool_set_true:N #1 }
}

\PrgNewFunction \BoolSetFalse { M }
{
    __fun_do_assignment:Nnn #1 { \bool_gset_false:N #1 } { \bool_set_false:N #1 }
}

```

```

\PrGNewFunction \BoolSetEq { M M }
{
  \__fun_do_assignment:Nnn #1
  { \bool_gset_eq:NN #1 #2 } { \bool_set_eq:NN #1 #2 }
}

\PrGNewFunction \BoolLog { m } { \bool_log:n { #1 } }

\PrGNewFunction \BoolVarLog { M } { \bool_log:N #1 }

\PrGNewFunction \BoolShow { m } { \bool_show:n { #1 } }

\PrGNewFunction \BoolVarShow { M } { \bool_show:N #1 }

\PrGNewConditional \BoolIfExist { M }
{
  \bool_if_exist:NTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \BoolVarIf { M } { \Result { #1 } }

\PrGNewConditional \BoolVarNot { M }
{
  \bool_if:NTF #1
  { \Result { \cFalseBool } } { \Result { \cTrueBool } }
}

\PrGNewConditional \BoolVarAnd { M M }
{
  \bool_lazy_and:nnTF {#1} {#2}
  { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \BoolVarOr { M M }
{
  \bool_lazy_or:nnTF {#1} {#2}
  { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \BoolVarXor { M M }
{
  \bool_xor:nnTF {#1} {#2}
  { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

```

11.4 Interfaces for Token Lists (Tl)

```

\tl_set_eq:NN \cEmptyTl \c_empty_tl
\tl_set_eq:NN \cSpaceTl \c_space_tl
\tl_set_eq:NN \cNoValueTl \c_novalue_tl

\tl_new:N \lTmptl \tl_new:N \lTmptl \tl_new:N \lTmptl
\tl_new:N \lTmptl \tl_new:N \lTmptl \tl_new:N \lTmptl
\tl_new:N \l@Funx@Tl \tl_new:N \l@Funy@Tl \tl_new:N \l@Funz@Tl

```

```

\tl_new:N \gTmptl    \tl_new:N \gTmptbTl    \tl_new:N \gTmptcTl
\tl_new:N \gTmptiTl    \tl_new:N \gTmptjTl    \tl_new:N \gTmptkTl
\tl_new:N \g@Funx@Tl    \tl_new:N \g@Funy@Tl    \tl_new:N \g@Funz@Tl

\PrgNewFunction \TlNew { M } { \tl_new:N #1 }

\PrgNewFunction \TlLog { m } { \tl_log:n { #1 } }

\PrgNewFunction \TlVarLog { M } { \tl_log:N #1 }

\PrgNewFunction \TlShow { m } { \tl_show:n { #1 } }

\PrgNewFunction \TlVarShow { M } { \tl_show:N #1 }

\PrgNewFunction \TlUse { M } { \Result { \Value #1 } }

\PrgNewFunction \TlToStr { m } { \Expand { \tl_to_str:n { #1 } } }

\PrgNewFunction \TlVarToStr { M } { \Expand { \tl_to_str:N #1 } }

\PrgNewFunction \TlConst { M m } { \tl_const:Nn #1 { #2 } }

\PrgNewFunction \TlSet { M m }
{
  \__fun_do_assignment:Nnn #1 { \tl_gset:Nn #1 {#2} } { \tl_set:Nn #1 {#2} }
}

\PrgNewFunction \TlSetEq { M M }
{
  \__fun_do_assignment:Nnn #1 { \tl_gset_eq:NN #1 #2 } { \tl_set_eq:NN #1 #2 }
}

\PrgNewFunction \TlConcat { M M M }
{
  \__fun_do_assignment:Nnn #1
  { \tl_gconcat:NNN #1 #2 #3 } { \tl_concat:NNN #1 #2 #3 }
}

\PrgNewFunction \TlClear { M }
{
  \__fun_do_assignment:Nnn #1 { \tl_gclear:N #1 } { \tl_clear:N #1 }
}

\PrgNewFunction \TlClearNew { M }
{
  \__fun_do_assignment:Nnn #1 { \tl_gclear_new:N #1 } { \tl_clear_new:N #1 }
}

\PrgNewFunction \TlPutLeft { M m }
{
  \__fun_do_assignment:Nnn #1
  { \tl_gput_left:Nn #1 {#2} } { \tl_put_left:Nn #1 {#2} }
}

```

```

\PrgNewFunction \TlPutRight { M m }
{
  \__fun_do_assignment:Nnn #1
  { \tl_gput_right:Nn #1 {#2} } { \tl_put_right:Nn #1 {#2} }
}

\PrgNewFunction \TlReplaceOnce { M m m }
{
  \__fun_do_assignment:Nnn #1
  { \tl_greplace_once:Nnn #1 {#2} {#3} } { \tl_replace_once:Nnn #1 {#2} {#3} }
}

\PrgNewFunction \TlReplaceAll { M m m }
{
  \__fun_do_assignment:Nnn #1
  { \tl_greplace_all:Nnn #1 {#2} {#3} } { \tl_replace_all:Nnn #1 {#2} {#3} }
}

\PrgNewFunction \TlRemoveOnce { M m }
{
  \__fun_do_assignment:Nnn #1
  { \tl_gremove_once:Nn #1 {#2} } { \tl_remove_once:Nn #1 {#2} }
}

\PrgNewFunction \TlRemoveAll { M m }
{
  \__fun_do_assignment:Nnn #1
  { \tl_gremove_all:Nn #1 {#2} } { \tl_remove_all:Nn #1 {#2} }
}

\PrgNewFunction \TlTrimSpaces { m } { \Expand { \tl_trim_spaces:n { #1 } } }

\PrgNewFunction \TlVarTrimSpaces { M }
{
  \__fun_do_assignment:Nnn #1 { \tl_gtrim_spaces:N #1 } { \tl_trim_spaces:N #1 }
}

\PrgNewFunction \TlCount { m } { \Expand { \tl_count:n { #1 } } }

\PrgNewFunction \TlVarCount { M } { \Expand { \tl_count:N #1 } }

\PrgNewFunction \TlHead { m } { \Expand { \tl_head:n { #1 } } }

\PrgNewFunction \TlVarHead { M } { \Expand { \tl_head:N #1 } }

\PrgNewFunction \TlTail { m } { \Expand { \tl_tail:n { #1 } } }

\PrgNewFunction \TlVarTail { M } { \Expand { \tl_tail:N #1 } }

\PrgNewFunction \TlItem { m m } { \Expand { \tl_item:nn {#1} {#2} } }

\PrgNewFunction \TlVarItem { M m } { \Expand { \tl_item:Nn #1 {#2} } }

\PrgNewFunction \TlRandItem { m } { \Expand { \tl_rand_item:n {#1} } }

```

```

\PrGNewFunction \TlVarRandItem { M } { \Expand { \tl_rand_item:N #1 } }

\PrGNewFunction \TlVarCase { M m } { \tl_case:Nn {#1} {#2} }
\PrGNewFunction \TlVarCaseT { M m n } { \tl_case:NnT {#1} {#2} {#3} }
\PrGNewFunction \TlVarCaseF { M m n } { \tl_case:NnF {#1} {#2} {#3} }
\PrGNewFunction \TlVarCaseTF { M m n n } { \tl_case:NnTF {#1} {#2} {#3} {#4} }

\PrGNewFunction \TlMapInline { m n }
{
  \tl_map_inline:nn {#1} {#2}
}

\PrGNewFunction \TlVarMapInline { M n }
{
  \tl_map_inline:Nn #1 {#2}
}

\PrGNewFunction \TlMapVariable { m M n }
{
  \tl_map_variable:nNn {#1} #2 {#3}
}

\PrGNewFunction \TlVarMapVariable { M M n }
{
  \tl_map_variable:NNn #1 #2 {#3}
}

\PrGNewConditional \TlIfExist { M }
{
  \tl_if_exist:NnTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \TlIfEmpty { m }
{
  \tl_if_empty:nTF {#1} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \TlVarIfEmpty { M }
{
  \tl_if_empty:NnTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \TlIfBlank { m }
{
  \tl_if_blank:nTF {#1} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \TlIfEq { m m }
{
  \tl_if_eq:nnTF {#1} {#2} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \TlVarIfEq { M M }
{
  \tl_if_eq:NnTF #1 #2 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

```

```

}

\PrGNewConditional \TlIfIn { m m }
{
  \tl_if_in:nnTF {#1} {#2} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \TlVarIfIn { M m }
{
  \tl_if_in:NnTF #1 {#2} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \TlIfSingle { m }
{
  \tl_if_single:nTF {#1} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \TlVarIfSingle { M }
{
  \tl_if_single:NTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

```

11.5 Interfaces for Strings (Str)

```

\str_set_eq:NN \cAmpersandStr \c_ampersand_str
\str_set_eq:NN \cAttignStr \c_atign_str
\str_set_eq:NN \cBackslashStr \c_backslash_str
\str_set_eq:NN \cLeftBraceStr \c_left_brace_str
\str_set_eq:NN \cRightBraceStr \c_right_brace_str
\str_set_eq:NN \cCircumflexStr \c_circumflex_str
\str_set_eq:NN \cColonStr \c_colon_str
\str_set_eq:NN \cDollarStr \c_dollar_str
\str_set_eq:NN \cHashStr \c_hash_str
\str_set_eq:NN \cPercentStr \c_percent_str
\str_set_eq:NN \cTildeStr \c_tilde_str
\str_set_eq:NN \cUnderscoreStr \c_underscore_str
\str_set_eq:NN \cZeroStr \c_zero_str

\str_new:N \lTmPaStr \str_new:N \lTmPbStr \str_new:N \lTmPcStr
\str_new:N \lTmPiStr \str_new:N \lTmPjStr \str_new:N \lTmPkStr
\str_new:N \l@Funx@Str \str_new:N \l@Funy@Str \str_new:N \l@Funz@Str

\str_new:N \gTmPaStr \str_new:N \gTmPbStr \str_new:N \gTmPcStr
\str_new:N \gTmPiStr \str_new:N \gTmPjStr \str_new:N \gTmPkStr
\str_new:N \g@Funx@Str \str_new:N \g@Funy@Str \str_new:N \g@Funz@Str

\PrGNewFunction \StrNew { M } { \str_new:N #1 }

\PrGNewFunction \StrLog { m } { \str_log:n { #1 } }

\PrGNewFunction \StrVarLog { M } { \str_log:N #1 }

\PrGNewFunction \StrShow { m } { \str_show:n { #1 } }

```

```

\PrGNewFunction \StrVarShow { M } { \str_show:N #1 }

\PrGNewFunction \StrUse { M } { \Result { \Value #1 } }

\PrGNewFunction \StrConst { M m } { \str_const:Nn #1 {#2} }

\PrGNewFunction \StrSet { M m }
{
  \__fun_do_assignment:Nnn #1 { \str_gset:Nn #1 {#2} } { \str_set:Nn #1 {#2} }
}

\PrGNewFunction \StrSetEq { M M }
{
  \__fun_do_assignment:Nnn #1 { \str_gset_eq:NN #1 #2 } { \str_set_eq:NN #1 #2 }
}

\PrGNewFunction \StrConcat { M M M }
{
  \__fun_do_assignment:Nnn #1
  { \str_gconcat:NNN #1 #2 #3 } { \str_concat:NNN #1 #2 #3 }
}

\PrGNewFunction \StrClear { M }
{
  \__fun_do_assignment:Nnn #1 { \str_gclear:N #1 } { \str_clear:N #1 }
}

\PrGNewFunction \StrClearNew { M }
{
  \__fun_do_assignment:Nnn #1 { \str_gclear_new:N #1 } { \str_clear_new:N #1 }
}

\PrGNewFunction \StrPutLeft { M m }
{
  \__fun_do_assignment:Nnn #1
  { \str_gput_left:Nn #1 {#2} } { \str_put_left:Nn #1 {#2} }
}

\PrGNewFunction \StrPutRight { M m }
{
  \__fun_do_assignment:Nnn #1
  { \str_gput_right:Nn #1 {#2} } { \str_put_right:Nn #1 {#2} }
}

\PrGNewFunction \StrReplaceOnce { M m m }
{
  \__fun_do_assignment:Nnn #1
  { \str_greplace_once:Nnn #1 {#2} {#3} } { \str_replace_once:Nnn #1 {#2} {#3} }
}

\PrGNewFunction \StrReplaceAll { M m m }
{
  \__fun_do_assignment:Nnn #1
  { \str_greplace_all:Nnn #1 {#2} {#3} } { \str_replace_all:Nnn #1 {#2} {#3} }
}

```

```

\PrGNewFunction \StrRemoveOnce { M m }
{
  \__fun_do_assignment:Nnn #1
  { \str_gremove_once:Nn #1 {#2} } { \str_remove_once:Nn #1 {#2} }
}

\PrGNewFunction \StrRemoveAll { M m }
{
  \__fun_do_assignment:Nnn #1
  { \str_gremove_all:Nn #1 {#2} } { \str_remove_all:Nn #1 {#2} }
}

%% Avoid naming conflict with xstring package
\cs_if_exist:NF \StrCount
{ \PrGNewFunction \StrCount { m } { \Expand { \str_count:n { #1 } } } }

%% Provide another name for \StrCount function
\PrGNewFunction \StrSize { m } { \Expand { \str_count:n { #1 } } }

\PrGNewFunction \StrVarCount { M } { \Expand { \str_count:N #1 } }

\PrGNewFunction \StrHead { m } { \Expand { \str_head:n { #1 } } }

\PrGNewFunction \StrVarHead { M } { \Expand { \str_head:N #1 } }

\PrGNewFunction \StrTail { m } { \Expand { \str_tail:n { #1 } } }

\PrGNewFunction \StrVarTail { M } { \Expand { \str_tail:N #1 } }

\PrGNewFunction \StrItem { m m } { \Expand { \str_item:nn {#1} {#2} } }

\PrGNewFunction \StrVarItem { M m } { \Expand { \str_item:Nn #1 {#2} } }

\PrGNewFunction \StrCase { m m } { \str_case:nn {#1} {#2} }
\PrGNewFunction \StrCaseT { m m n } { \str_case:nnT {#1} {#2} {#3} }
\PrGNewFunction \StrCaseF { m m n } { \str_case:nnF {#1} {#2} {#3} }
\PrGNewFunction \StrCaseTF { m m n n } { \str_case:nnTF {#1} {#2} {#3} {#4} }

\PrGNewFunction \StrMapInline { m n }
{
  \str_map_inline:nn {#1} {#2}
}

\PrGNewFunction \StrVarMapInline { M n }
{
  \str_map_inline:Nn #1 {#2}
}

\PrGNewFunction \StrMapVariable { m M n }
{
  \str_map_variable:nNn {#1} #2 {#3}
}

\PrGNewFunction \StrVarMapVariable { M M n }

```



```

{
    \str_map_variable:NNn #1 #2 {#3}
}

\PrgNewConditional \StrIfExist { M }
{
    \str_if_exist:NTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrgNewConditional \StrVarIfEmpty { M }
{
    \str_if_empty:NTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrgNewConditional \StrIfEq { m m }
{
    \str_if_eq:nnTF {#1} {#2} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrgNewConditional \StrVarIfEq { M M }
{
    \str_if_eq:NNTF #1 #2 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrgNewConditional \StrIfIn { m m }
{
    \str_if_in:nnTF {#1} {#2} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrgNewConditional \StrVarIfIn { M m }
{
    \str_if_in:NnTF #1 {#2} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

%% Avoid naming conflict with xstring package
\cs_if_exist:NF \StrCompare
{
    \PrgNewConditional \StrCompare { m N m }
    {
        \str_compare:nNnTF {#1} #2 {#3}
        { \Result { \cTrueBool } }
        { \Result { \cFalseBool } }
    }
}

%% Provide another name for \StrCompare function
\PrgNewConditional \StrIfCompare { m N m }
{
    \str_compare:nNnTF {#1} #2 {#3}
    { \Result { \cTrueBool } }
    { \Result { \cFalseBool } }
}

```

11.6 Interfaces for Integers (Int)

```

\cs_set_eq:NN \cZeroInt          \c_zero_int

```

```

\cs_set_eq:NN \cOneInt      \c_one_int
\cs_set_eq:NN \cMaxInt      \c_max_int
\cs_set_eq:NN \cMaxRegisterInt \c_max_register_int
\cs_set_eq:NN \cMaxCharInt   \c_max_char_in

\int_new:N \lTmptaInt   \int_new:N \lTmpbInt   \int_new:N \lTmpcInt
\int_new:N \lTmpiInt    \int_new:N \lTmpjInt    \int_new:N \lTmpkInt
\int_new:N \l@Funx@Int  \int_new:N \l@Funy@Int  \int_new:N \l@Funz@Int

\int_new:N \gTmptaInt   \int_new:N \gTmpbInt   \int_new:N \gTmpcInt
\int_new:N \gTmpiInt    \int_new:N \gTmpjInt    \int_new:N \gTmpkInt
\int_new:N \g@Funx@Int  \int_new:N \g@Funy@Int  \int_new:N \g@Funz@Int

\PrgNewFunction \IntEval { m }
{
  \Result { \Expand { \int_eval:n { #1 } } }
}

\PrgNewFunction \IntMathAdd { m m }
{
  \int_set:Nn \l@Funx@Int { \int_eval:n { (#1) + (#2) } }
  \Result { \Value \l@Funx@Int }
}

\PrgNewFunction \IntMathSub { m m }
{
  \int_set:Nn \l@Funx@Int { \int_eval:n { (#1) - (#2) } }
  \Result { \Value \l@Funx@Int }
}

\PrgNewFunction \IntMathMult { m m }
{
  \int_set:Nn \l@Funx@Int { \int_eval:n { (#1) * (#2) } }
  \Result { \Value \l@Funx@Int }
}

\PrgNewFunction \IntMathDiv { m m }
{
  \Expand { \int_div_round:nn { #1 } { #2 } }
}

\PrgNewFunction \IntMathDivTruncate { m m }
{
  \Expand { \int_div_truncate:nn { #1 } { #2 } }
}

\PrgNewFunction \IntMathSign { m } { \Expand { \int_sign:n { #1 } } }

\PrgNewFunction \IntMathAbs { m } { \Expand { \int_abs:n { #1 } } }

\PrgNewFunction \IntMathMax { m m } { \Expand { \int_max:nn { #1 } { #2 } } }

\PrgNewFunction \IntMathMin { m m } { \Expand { \int_min:nn { #1 } { #2 } } }

\PrgNewFunction \IntMathMod { m m } { \Expand { \int_mod:nn { #1 } { #2 } } }

```

```

\PrGNewFunction \IntMathRand { m m } { \Expand { \int_rand:nn { #1 } { #2 } } } }

\PrGNewFunction \IntNew { M } { \int_new:N #1 }

\PrGNewFunction \IntConst { M m } { \int_const:Nn #1 { #2 } }

\PrGNewFunction \IntLog { m } { \int_log:n { #1 } }

\PrGNewFunction \IntVarLog { M } { \int_log:N #1 }

\PrGNewFunction \IntShow { m } { \int_show:n { #1 } }

\PrGNewFunction \IntVarShow { M } { \int_show:N #1 }

\PrGNewFunction \IntUse { M } { \Result { \Value #1 } }

\PrGNewFunction \IntSet { M m }
{
  \__fun_do_assignment:Nnn #1 { \int_gset:Nn #1 {#2} } { \int_set:Nn #1 {#2} }
}

\PrGNewFunction \IntZero { M }
{
  \__fun_do_assignment:Nnn #1 { \int_gzero:N #1 } { \int_zero:N #1 }
}

\PrGNewFunction \IntZeroNew { M }
{
  \__fun_do_assignment:Nnn #1 { \int_gzero_new:N #1 } { \int_zero_new:N #1 }
}

\PrGNewFunction \IntSetEq { M M }
{
  \__fun_do_assignment:Nnn #1 { \int_gset_eq:NN #1 #2 } { \int_set_eq:NN #1 #2 }
}

\PrGNewFunction \IntIncr { M }
{
  \__fun_do_assignment:Nnn #1 { \int_gincr:N #1 } { \int_incr:N #1 }
}

\PrGNewFunction \IntDecr { M }
{
  \__fun_do_assignment:Nnn #1 { \int_gdecr:N #1 } { \int_decr:N #1 }
}

\PrGNewFunction \IntAdd { M m }
{
  \__fun_do_assignment:Nnn #1 { \int_gadd:Nn #1 {#2} } { \int_add:Nn #1 {#2} }
}

\PrGNewFunction \IntSub { M m }
{
  \__fun_do_assignment:Nnn #1 { \int_gsub:Nn #1 {#2} } { \int_sub:Nn #1 {#2} }
}

```

```

}

\PrGNewFunction \IntStepInline { m m m n }
{
  \int_step_inline:nnnn { #1 } { #2 } { #3 } { #4 }
}

\PrGNewFunction \IntStepVariable { m m m M n }
{
  \int_step_variable:nnnNn { #1 } { #2 } { #3 } #4 { #5 }
}

\PrGNewConditional \IntIfExist { M }
{
  \int_if_exist:NTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \IntIfOdd { m }
{
  \int_if_odd:nTF { #1 } { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \IntIfEven { m }
{
  \int_if_even:nTF { #1 } { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \IntCompare { m N m }
{
  \int_compare:nNnTF {#1} #2 {#3}
  { \Result { \cTrueBool } }
  { \Result { \cFalseBool } }
}

\PrGNewFunction \IntCase { m m } { \int_case:nn {#1} {#2} }
\PrGNewFunction \IntCaseT { m m n } { \int_case:nnT {#1} {#2} {#3} }
\PrGNewFunction \IntCaseF { m m n } { \int_case:nnF {#1} {#2} {#3} }
\PrGNewFunction \IntCaseTF { m m n n } { \int_case:nnTF {#1} {#2} {#3} {#4} }

```

11.7 Interfaces for Floating Point Numbers (Fp)

```

\fp_set_eq:NN \cZeroFp      \c_zero_fp
\fp_set_eq:NN \cMinusZeroFp \c_minus_zero_fp
\fp_set_eq:NN \cOneFp       \c_one_fp
\fp_set_eq:NN \cInfFp       \c_inf_fp
\fp_set_eq:NN \cMinusInfFp  \c_minus_inf_fp
\fp_set_eq:NN \cEFp         \c_e_fp
\fp_set_eq:NN \cPiFp        \c_pi_fp
\fp_set_eq:NN \cOneDegreeFp \c_one_degree_fp

\fp_new:N \lTmPaFp   \fp_new:N \lTmPbFp   \fp_new:N \lTmPcFp
\fp_new:N \lTmPiFp   \fp_new:N \lTmPjFp   \fp_new:N \lTmPkFp
\fp_new:N \l@Funx@Fp \fp_new:N \l@Funy@Fp \fp_new:N \l@Funz@Fp

\fp_new:N \gTmPaFp   \fp_new:N \gTmPbFp   \fp_new:N \gTmPcFp

```

```

\fp_new:N \gTmpiFp   \fp_new:N \gTmpjFp   \fp_new:N \gTmpkFp
\fp_new:N \g@Funx@Fp \fp_new:N \g@Funy@Fp \fp_new:N \g@Funz@Fp

```

```

\PrgNewFunction \FpEval { m }
{
  \Result { \Expand { \fp_eval:n { #1 } } }
}

```

```

\PrgNewFunction \FpMathAdd { m m }
{
  \fp_set:Nn \l@Funx@Fp { \fp_eval:n { (#1) + (#2) } }
  \Result { \FpUse \l@Funx@Fp }
}

```

```

\PrgNewFunction \FpMathSub { m m }
{
  \fp_set:Nn \l@Funx@Fp { \fp_eval:n { (#1) - (#2) } }
  \Result { \FpUse \l@Funx@Fp }
}

```

```

\PrgNewFunction \FpMathMult { m m }
{
  \fp_set:Nn \l@Funx@Fp { \fp_eval:n { (#1) * (#2) } }
  \Result { \FpUse \l@Funx@Fp }
}

```

```

\PrgNewFunction \FpMathDiv { m m }
{
  \fp_set:Nn \l@Funx@Fp { \fp_eval:n { (#1) / (#2) } }
  \Result { \FpUse \l@Funx@Fp }
}

```

```

\PrgNewFunction \FpMathSign { m }
{
  \Result { \Expand { \fp_sign:n { #1 } } }
}

```

```

\PrgNewFunction \FpMathAbs { m }
{
  \Result { \Expand { \fp_abs:n { #1 } } }
}

```

```

\PrgNewFunction \FpMathMax { m m }
{
  \Result { \Expand { \fp_max:nn { #1 } { #2 } } }
}

```

```

\PrgNewFunction \FpMathMin { m m }
{
  \Result { \Expand { \fp_min:nn { #1 } { #2 } } }
}

```

```

\PrgNewFunction \FpNew { M } { \fp_new:N #1 }

```

```

\PrgNewFunction \FpConst { M m } { \fp_const:Nn #1 {#2} }

```

```

\PrNewFunction \FpUse { M } { \Result { \Expand { \fp_use:N #1 } } }

\PrNewFunction \FpLog { m } { \fp_log:n { #1 } }

\PrNewFunction \FpVarLog { M } { \fp_log:N #1 }

\PrNewFunction \FpShow { m } { \fp_show:n { #1 } }

\PrNewFunction \FpVarShow { M } { \fp_show:N #1 }

\PrNewFunction \FpSet { M m }
{
  \__fun_do_assignment:Nnn #1 { \fp_gset:Nn #1 {#2} } { \fp_set:Nn #1 {#2} }
}

\PrNewFunction \FpSetEq { M M }
{
  \__fun_do_assignment:Nnn #1 { \fp_gset_eq:NN #1 #2 } { \fp_set_eq:NN #1 #2 }
}

\PrNewFunction \FpZero { M }
{
  \__fun_do_assignment:Nnn #1 { \fp_gzero:N #1 } { \fp_zero:N #1 }
}

\PrNewFunction \FpZeroNew { M }
{
  \__fun_do_assignment:Nnn #1 { \fp_gzero_new:N #1 } { \fp_zero_new:N #1 }
}

\PrNewFunction \FpAdd { M m }
{
  \__fun_do_assignment:Nnn #1 { \fp_gadd:Nn #1 {#2} } { \fp_add:Nn #1 {#2} }
}

\PrNewFunction \FpSub { M m }
{
  \__fun_do_assignment:Nnn #1 { \fp_gsub:Nn #1 {#2} } { \fp_sub:Nn #1 {#2} }
}

\PrNewFunction \FpStepInline { m m m n }
{
  \fp_step_inline:nnnn { #1 } { #2 } { #3 } { #4 }
}

\PrNewFunction \FpStepVariable { m m m M n }
{
  \fp_step_variable:nnnNn { #1 } { #2 } { #3 } #4 { #5 }
}

\PrNewConditional \FpIfExist { M }
{
  \fp_if_exist:NTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

```

```

\PrGNewConditional \FpCompare { m N m }
{
  \fp_compare:nNnTF {#1} #2 {#3}
  { \Result { \cTrueBool } }
  { \Result { \cFalseBool } }
}

```

11.8 Interfaces for Dimensions (Dim)

```

\cs_set_eq:NN \cMaxDim \c_max_dim
\cs_set_eq:NN \cZeroDim \c_zero_dim

```

```

\dim_new:N \lTmPaDim \dim_new:N \lTmPbDim \dim_new:N \lTmPcDim
\dim_new:N \lTmPiDim \dim_new:N \lTmPjDim \dim_new:N \lTmPkDim
\dim_new:N \l@Funx@Dim \dim_new:N \l@Funy@Dim \dim_new:N \l@Funz@Dim

```

```

\dim_new:N \gTmPaDim \dim_new:N \gTmPbDim \dim_new:N \gTmPcDim
\dim_new:N \gTmPiDim \dim_new:N \gTmPjDim \dim_new:N \gTmPkDim
\dim_new:N \g@Funx@Dim \dim_new:N \g@Funy@Dim \dim_new:N \g@Funz@Dim

```

```

\PrGNewFunction \DimEval { m }
{
  \Result { \Expand { \dim_eval:n { #1 } } }
}

```

```

\PrGNewFunction \DimMathAdd { m m }
{
  \dim_set:Nn \l@Funx@Dim { \dim_eval:n { (#1) + (#2) } }
  \Result { \Value \l@Funx@Dim }
}

```

```

\PrGNewFunction \DimMathSub { m m }
{
  \dim_set:Nn \l@Funx@Dim { \dim_eval:n { (#1) - (#2) } }
  \Result { \Value \l@Funx@Dim }
}

```

```

\PrGNewFunction \DimMathSign { m }
{
  \Result { \Expand { \dim_sign:n { #1 } } }
}

```

```

\PrGNewFunction \DimMathAbs { m }
{
  \Result { \Expand { \dim_abs:n { #1 } } }
}

```

```

\PrGNewFunction \DimMathMax { m m }
{
  \Result { \Expand { \dim_max:nn { #1 } { #2 } } }
}

```

```

\PrGNewFunction \DimMathMin { m m }
{

```

```

    \Result { \Expand { \dim_min:nn { #1 } { #2 } } }
}

\PrgNewFunction \DimMathRatio { m m }
{
    \Result { \Expand { \dim_ratio:nn { #1 } { #2 } } }
}

\PrgNewFunction \DimNew { M } { \dim_new:N #1 }

\PrgNewFunction \DimConst { M m } { \dim_const:Nn #1 {#2} }

\PrgNewFunction \DimUse { M } { \Result { \Value #1 } }

\PrgNewFunction \DimLog { m } { \dim_log:n { #1 } }

\PrgNewFunction \DimVarLog { M } { \dim_log:N #1 }

\PrgNewFunction \DimShow { m } { \dim_show:n { #1 } }

\PrgNewFunction \DimVarShow { M } { \dim_show:N #1 }

\PrgNewFunction \DimSet { M m }
{
    \__fun_do_assignment:Nnn #1 { \dim_gset:Nn #1 {#2} } { \dim_set:Nn #1 {#2} }
}

\PrgNewFunction \DimSetEq { M M }
{
    \__fun_do_assignment:Nnn #1 { \dim_gset_eq:NN #1 #2 } { \dim_set_eq:NN #1 #2 }
}

\PrgNewFunction \DimZero { M }
{
    \__fun_do_assignment:Nnn #1 { \dim_gzero:N #1 } { \dim_zero:N #1 }
}

\PrgNewFunction \DimZeroNew { M }
{
    \__fun_do_assignment:Nnn #1 { \dim_gzero_new:N #1 } { \dim_zero_new:N #1 }
}

\PrgNewFunction \DimAdd { M m }
{
    \__fun_do_assignment:Nnn #1 { \dim_gadd:Nn #1 {#2} } { \dim_add:Nn #1 {#2} }
}

\PrgNewFunction \DimSub { M m }
{
    \__fun_do_assignment:Nnn #1 { \dim_gsub:Nn #1 {#2} } { \dim_sub:Nn #1 {#2} }
}

\PrgNewFunction \DimStepInline { m m m n }
{

```



```

    \dim_step_inline:nnnn { #1 } { #2 } { #3 } { #4 }
  }

\PrGNewFunction \DimStepVariable { m m m M n }
{
  \dim_step_variable:nnnNn { #1 } { #2 } { #3 } #4 { #5 }
}

\PrGNewConditional \DimIfExist { M }
{
  \dim_if_exist:NTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewConditional \DimCompare { m N m }
{
  \dim_compare:nNnTF {#1} #2 {#3}
    { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrGNewFunction \DimCase { m m } { \dim_case:nn {#1} {#2} }
\PrGNewFunction \DimCaseT { m m n } { \dim_case:nnT {#1} {#2} {#3} }
\PrGNewFunction \DimCaseF { m m n } { \dim_case:nnF {#1} {#2} {#3} }
\PrGNewFunction \DimCaseTF { m m n n } { \dim_case:nnTF {#1} {#2} {#3} {#4} }

```

11.9 Interfaces for Sorting Functions (Sort)

```

\cs_set_eq:NN \SortReturnSame \sort_return_same:
\cs_set_eq:NN \SortReturnSwapped \sort_return_swapped:

```

11.10 Interfaces for Comma Separated Lists (Clist)

```

\clist_new:N \lTmPaClist \clist_new:N \lTmPbClist \clist_new:N \lTmPcClist
\clist_new:N \lTmPiClist \clist_new:N \lTmPjClist \clist_new:N \lTmPkClist
\clist_new:N \l@Funx@Clist \clist_new:N \l@Funy@Clist \clist_new:N \l@Funz@Clist

\clist_new:N \gTmPaClist \clist_new:N \gTmPbClist \clist_new:N \gTmPcClist
\clist_new:N \gTmPiClist \clist_new:N \gTmPjClist \clist_new:N \gTmPkClist
\clist_new:N \g@Funx@Clist \clist_new:N \g@Funy@Clist \clist_new:N \g@Funz@Clist

\clist_set_eq:NN \cEmptyClist \c_empty_clist

\PrGNewFunction \ClistNew { M } { \clist_new:N #1 }

\PrGNewFunction \ClistLog { m } { \clist_log:n { #1 } }

\PrGNewFunction \ClistVarLog { M } { \clist_log:N #1 }

\PrGNewFunction \ClistShow { m } { \clist_show:n { #1 } }

\PrGNewFunction \ClistVarShow { M } { \clist_show:N #1 }

\PrGNewFunction \ClistVarJoin { M m }

```

```

{
  \Expand { \clist_use:Nn #1 { #2 } }
}

\PrgNewFunction \ClistVarJoinExtended { M m m m }
{
  \Expand { \clist_use:Nnn #1 { #2 } { #3 } { #4 } }
}

\PrgNewFunction \ClistJoin { m m }
{
  \Expand { \clist_use:nn { #1 } { #2 } }
}

\PrgNewFunction \ClistJoinExtended { m m m m }
{
  \Expand { \clist_use:nnnn { #1 } { #2 } { #3 } { #4 } }
}

\PrgNewFunction \ClistConst { M m } { \clist_const:Nn #1 { #2 } }

\PrgNewFunction \ClistSet { M m }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gset:Nn #1 {#2} } { \clist_set:Nn #1 {#2} }
}

\PrgNewFunction \ClistSetEq { M M }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gset_eq:NN #1 #2 } { \clist_set_eq:NN #1 #2 }
}

\PrgNewFunction \ClistSetFromSeq { M M }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gset_from_seq:NN #1 #2 } { \clist_set_from_seq:NN #1 #2 }
}

\PrgNewFunction \ClistConcat { M M M }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gconcat:NNN #1 #2 #3 } { \clist_concat:NNN #1 #2 #3 }
}

\PrgNewFunction \ClistClear { M }
{
  \__fun_do_assignment:Nnn #1 { \clist_gclear:N #1 } { \clist_clear:N #1 }
}

\PrgNewFunction \ClistClearNew { M }
{
  \__fun_do_assignment:Nnn #1 { \clist_gclear_new:N #1 } { \clist_clear_new:N #1 }
}

\PrgNewFunction \ClistPutLeft { M m }

```

```

{
  \__fun_do_assignment:Nnn #1
  { \clist_gput_left:Nn #1 {#2} } { \clist_put_left:Nn #1 {#2} }
}

\PrgNewFunction \ClistPutRight { M m }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gput_right:Nn #1 {#2} } { \clist_put_right:Nn #1 {#2} }
}

\PrgNewFunction \ClistRemoveDuplicates { M }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gremove_duplicates:N #1 } { \clist_remove_duplicates:N #1 }
}

\PrgNewFunction \ClistRemoveAll { M m }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gremove_all:Nn #1 {#2} } { \clist_remove_all:Nn #1 {#2} }
}

\PrgNewFunction \ClistReverse { M }
{
  \__fun_do_assignment:Nnn #1 { \clist_greverse:N #1 } { \clist_reverse:N #1 }
}

\PrgNewFunction \ClistSort { M m }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gsort:Nn #1 {#2} } { \clist_sort:Nn #1 {#2} }
}

\PrgNewFunction \ClistCount { m }
{
  \Result { \Expand { \clist_count:n { #1 } } }
}

\PrgNewFunction \ClistVarCount { m }
{
  \Result { \Expand { \clist_count:N #1 } }
}

\PrgNewFunction \ClistGet { M M } { \clist_get:NN #1 #2 }
\PrgNewFunction \ClistGetT { M M n } { \clist_get:NNT #1 #2 {#3} }
\PrgNewFunction \ClistGetF { M M n } { \clist_get:NNF #1 #2 {#3} }
\PrgNewFunction \ClistGetTF { M M n n } { \clist_get:NNTF #1 #2 {#3} {#4} }

\PrgNewFunction \ClistPop { M M }
{
  \__fun_do_assignment:Nnn #1
  { \clist_gpop:NN #1 #2 } { \clist_pop:NN #1 #2 }
}

\PrgNewFunction \ClistPopT { M M n }
{

```

```

    \__fun_do_assignment:Nnn #1
    { \clist_gpop:NNT #1 #2 {#3} } { \clist_pop:NNT #1 #2 {#3} }
}
\PrgNewFunction \ClistPopF { M M n }
{
    \__fun_do_assignment:Nnn #1
    { \clist_gpop:NNF #1 #2 {#3} } { \clist_pop:NNF #1 #2 {#3} }
}
\PrgNewFunction \ClistPopTF { M M n n }
{
    \__fun_do_assignment:Nnn #1
    { \clist_gpop:NNTF #1 #2 {#3} {#4} } { \clist_pop:NNTF #1 #2 {#3} {#4} }
}

\PrgNewFunction \ClistPush { M m }
{
    \__fun_do_assignment:Nnn #1
    { \clist_gpush:Nn #1 {#2} } { \clist_push:Nn #1 {#2} }
}

\PrgNewFunction \ClistItem { m m } { \Expand { \clist_item:nn {#1} {#2} } }

\PrgNewFunction \ClistVarItem { M m } { \Expand { \clist_item:Nn #1 {#2} } }

\PrgNewFunction \ClistRandItem { m } { \Expand { \clist_rand_item:n {#1} } }

\PrgNewFunction \ClistVarRandItem { M } { \Expand { \clist_rand_item:N #1 } }

\PrgNewFunction \ClistMapInline { m n }
{
    \clist_map_inline:nn {#1} {#2}
}

\PrgNewFunction \ClistVarMapInline { M n }
{
    \clist_map_inline:Nn #1 {#2}
}

\PrgNewFunction \ClistMapVariable { m M n }
{
    \clist_map_variable:nNn {#1} #2 {#3}
}

\PrgNewFunction \ClistVarMapVariable { M M n }
{
    \clist_map_variable:NNn #1 #2 {#3}
}

\cs_set_eq:NN \ClistMapBreak \clist_map_break:

\PrgNewConditional \ClistIfExist { M }
{
    \clist_if_exist:NTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

```

```

\PrNewConditional \ClistIfEmpty { m }
{
  \clist_if_empty:nTF {#1} { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrNewConditional \ClistVarIfEmpty { m }
{
  \clist_if_empty:NnTF #1 { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrNewConditional \ClistIfIn { m m }
{
  \clist_if_in:nnTF {#1} {#2}
  { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

\PrNewConditional \ClistVarIfIn { M m }
{
  \clist_if_in:NnTF #1 {#2}
  { \Result { \cTrueBool } } { \Result { \cFalseBool } }
}

```